A Tree Transducer Model for Synchronous Tree-Adjoining Grammars

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Synchronous Tree Substitution Grammar

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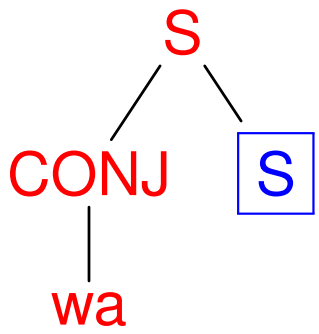
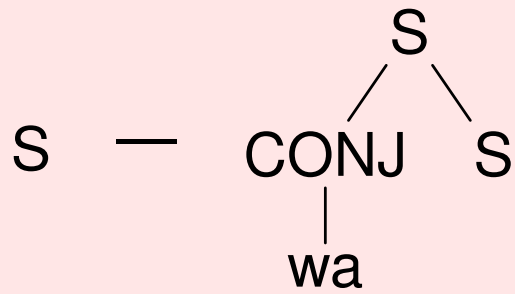


Used rule



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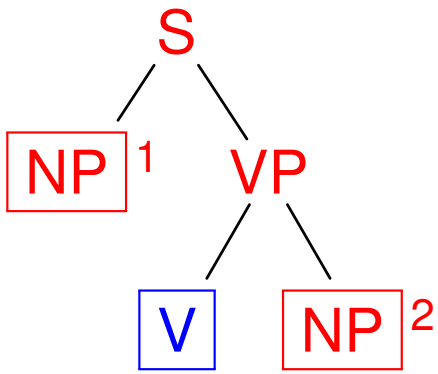

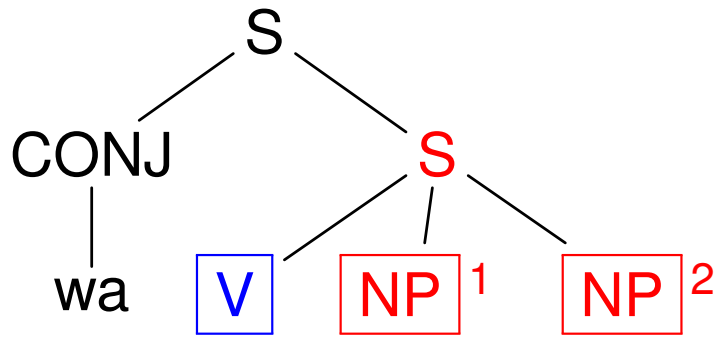

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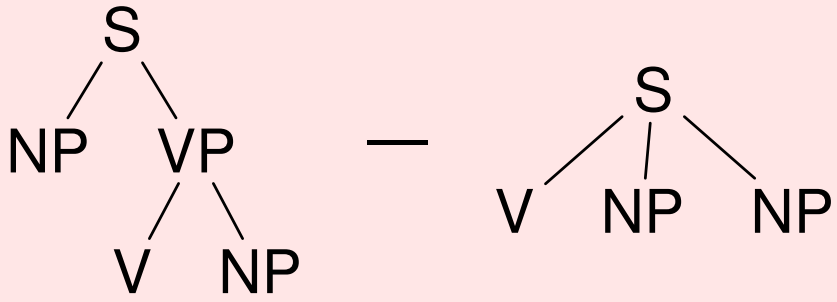

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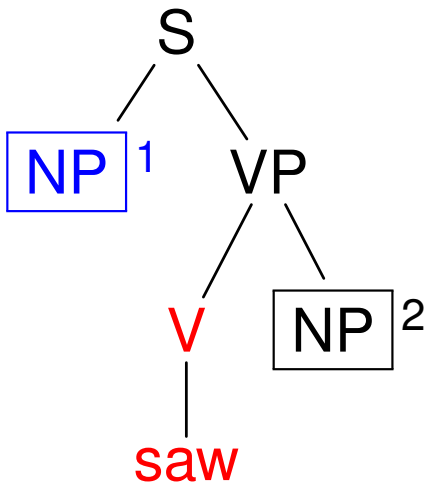

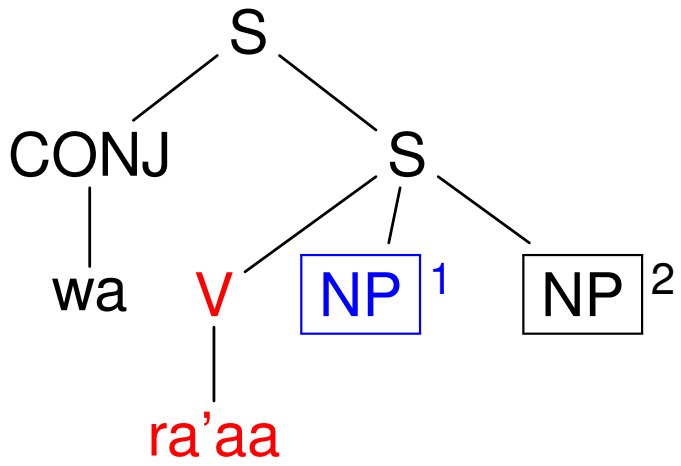

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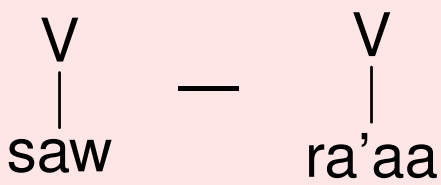

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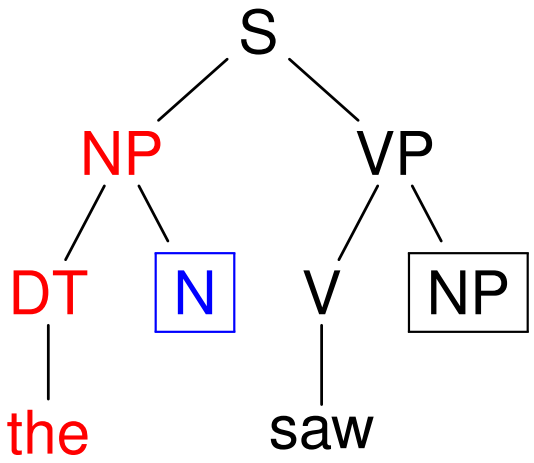

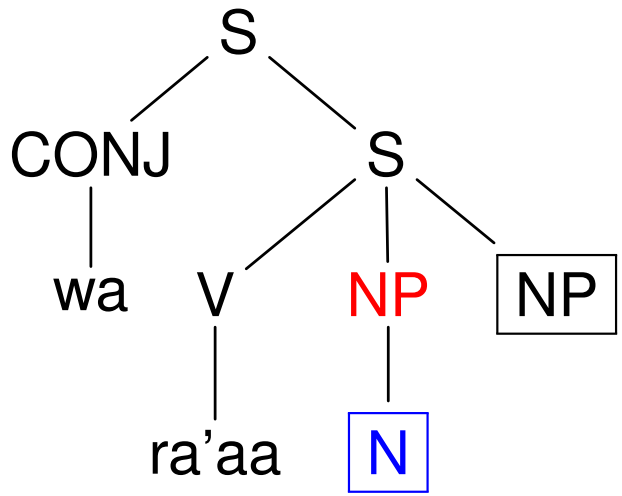

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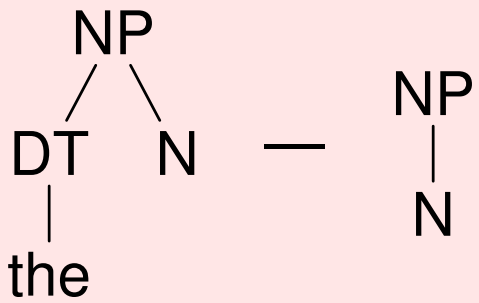

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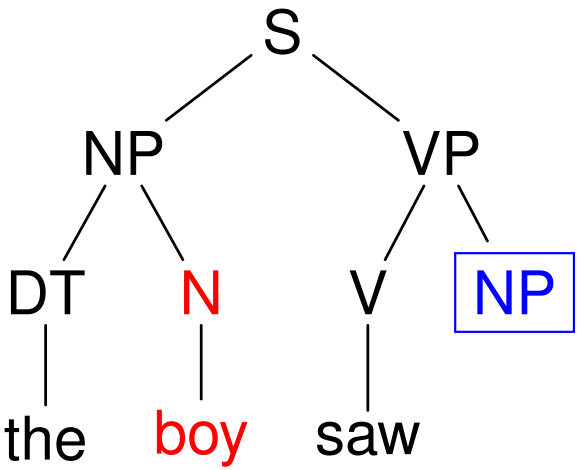

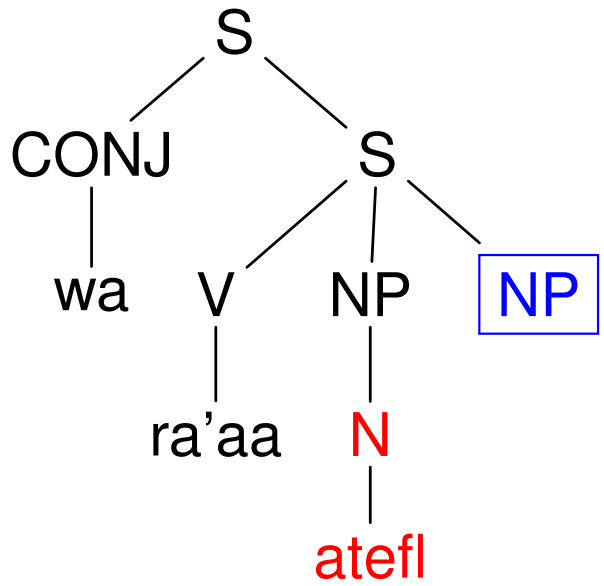

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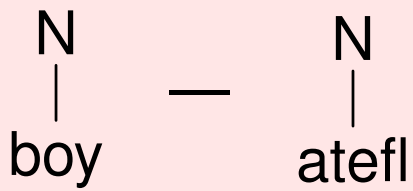

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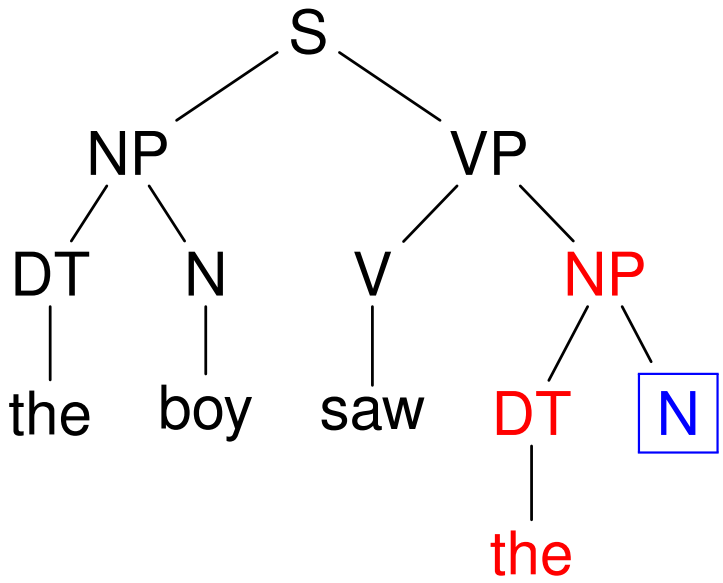

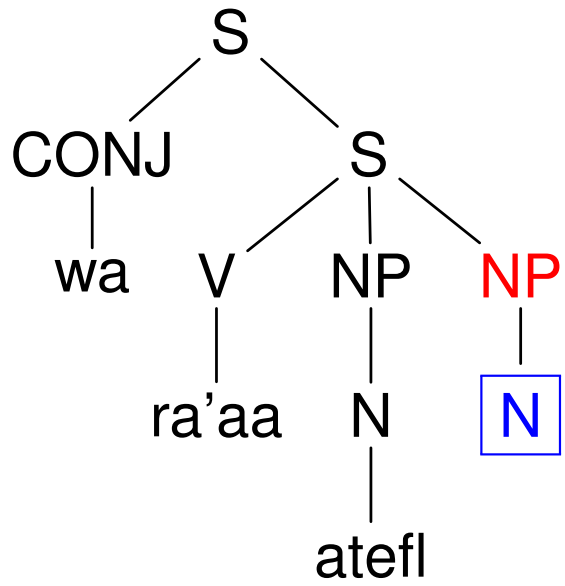

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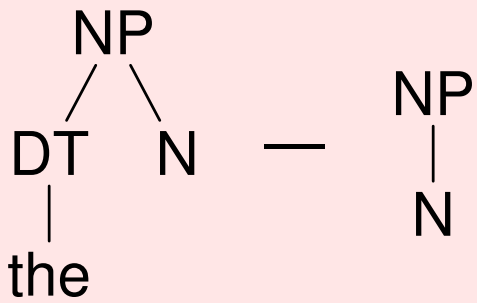

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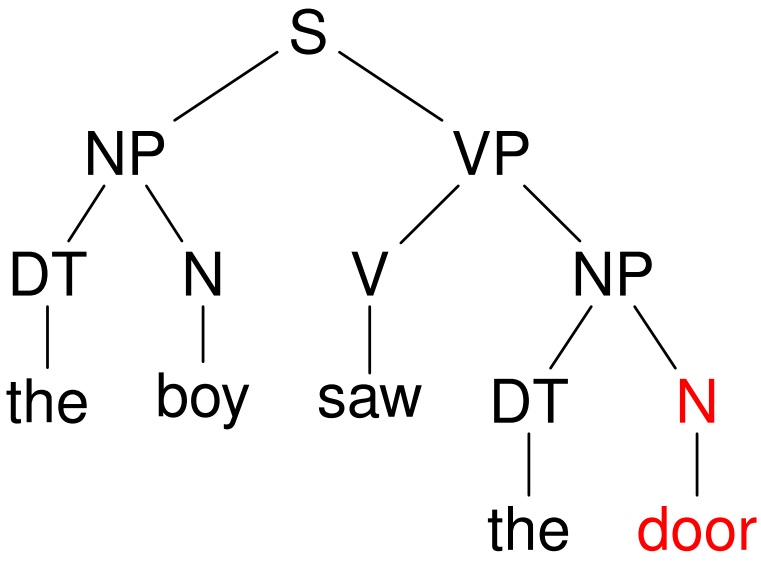
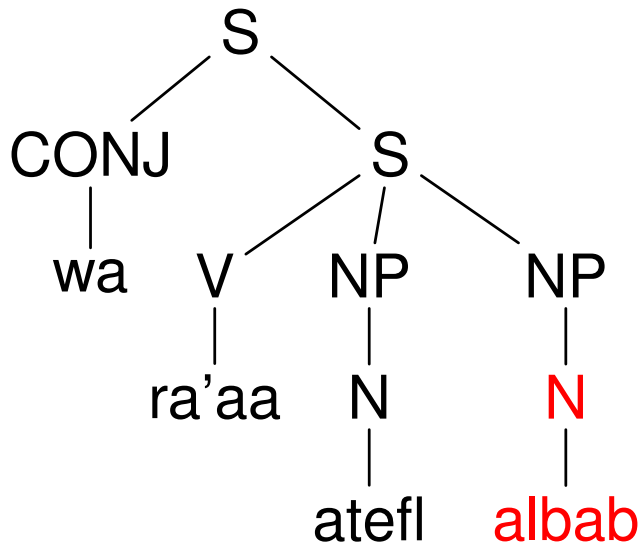
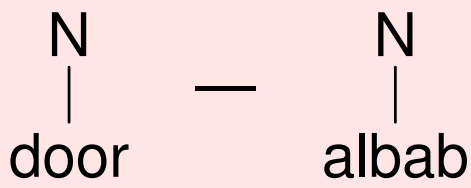




Used rule



Synchronous Tree Substitution Grammar (cont'd)

Advantages

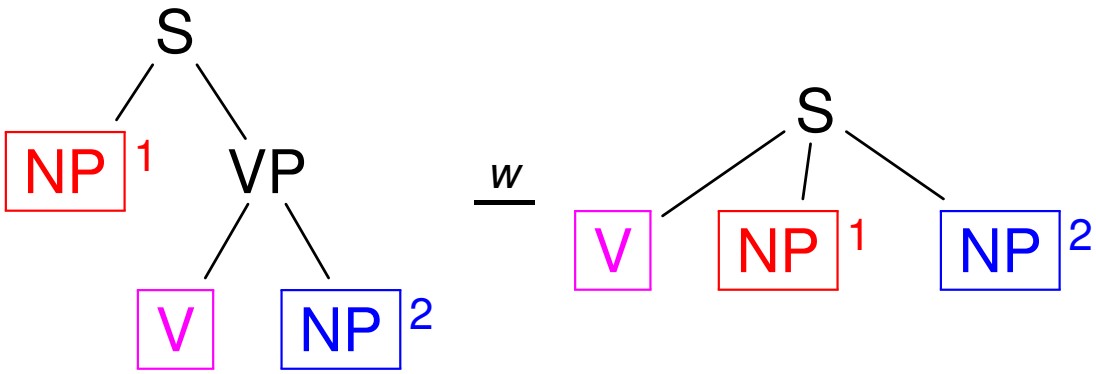
- simple and natural model
- easy to train (from linguistic resources)
- symmetric

Implementation

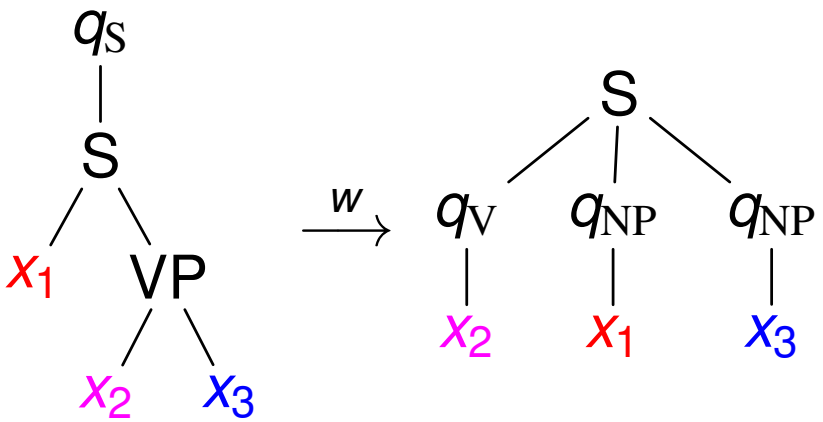
- extended top-down tree transducer in TIBURON [MAY, KNIGHT '06]

Synchronous Tree Substitution Grammar (cont'd)

Synchronous tree substitution grammar rule:



Corresponding extended top-down tree transducer rule:



Synchronous Tree-Adjoining Grammar

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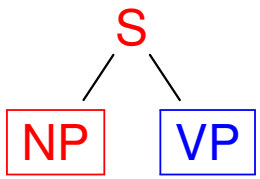
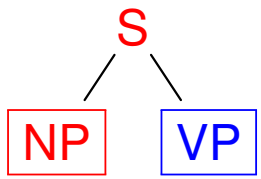
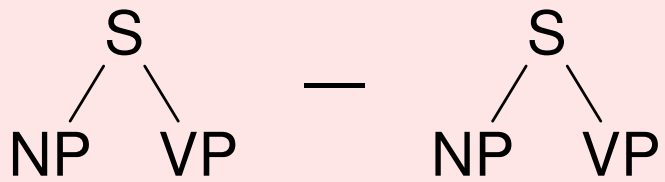




Used substitution rule



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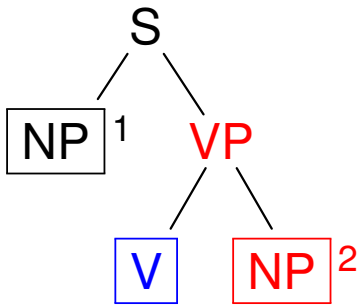

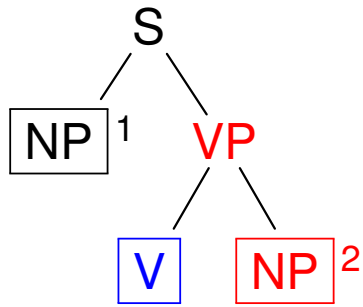

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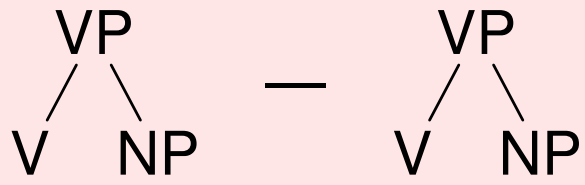

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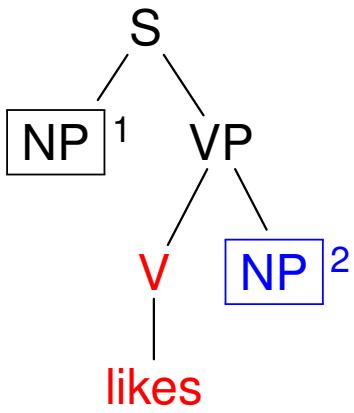

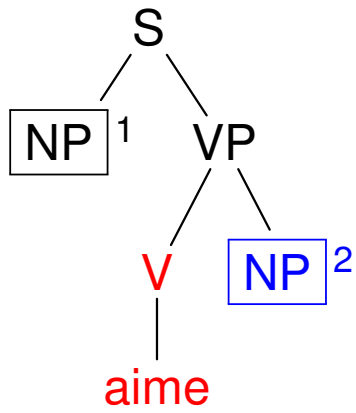

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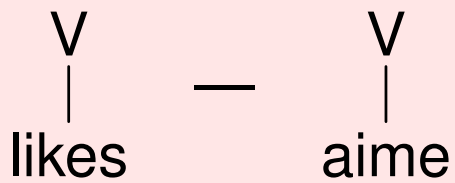

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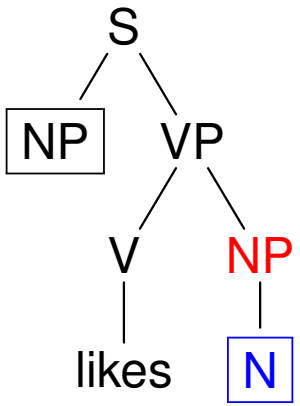

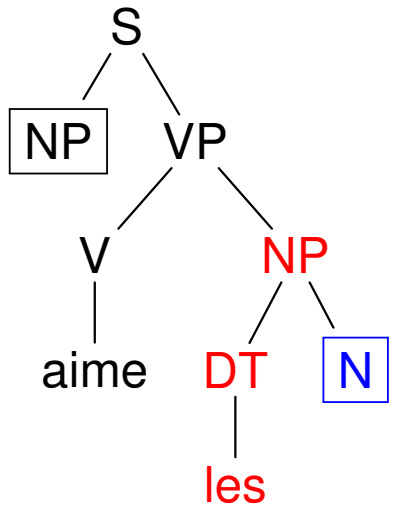

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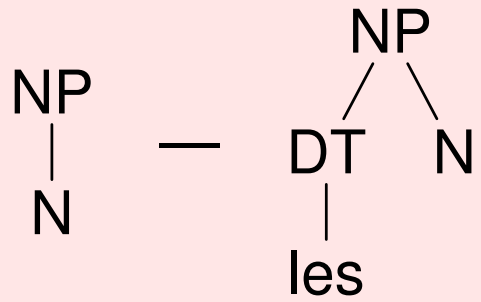

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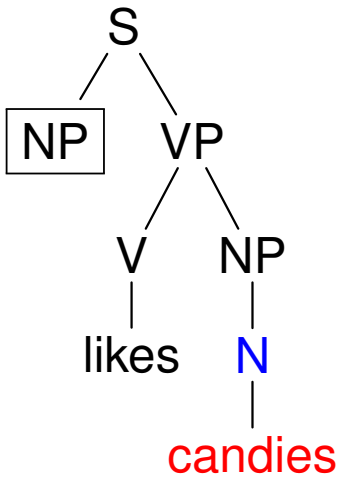

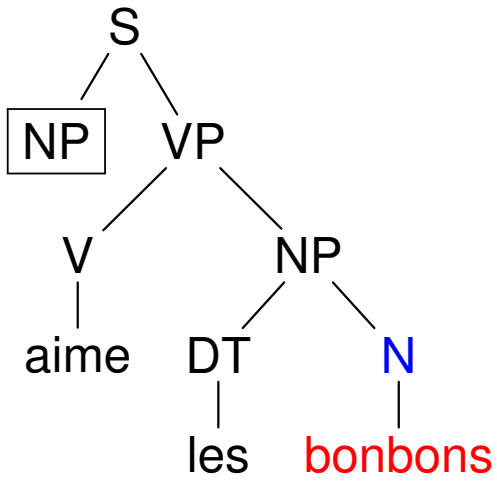

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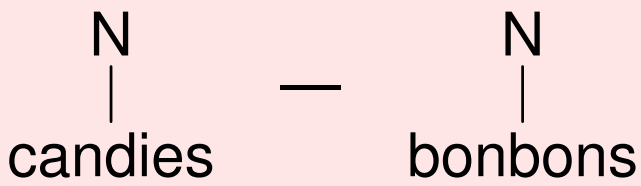

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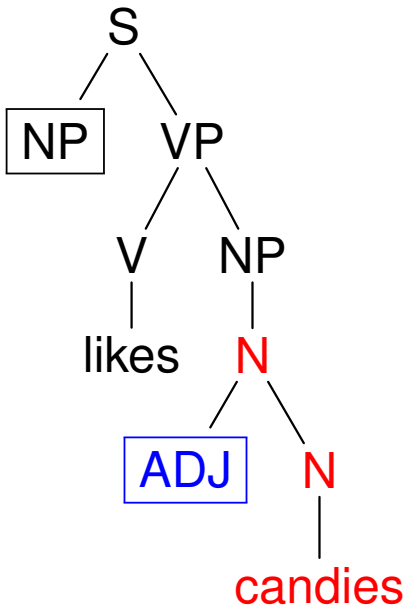
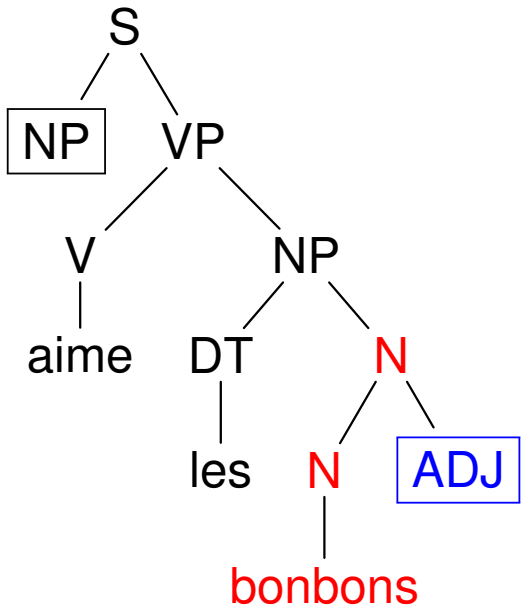
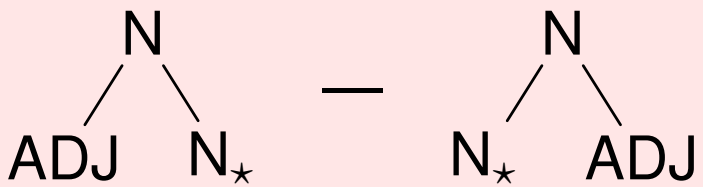




Used adjunction rule



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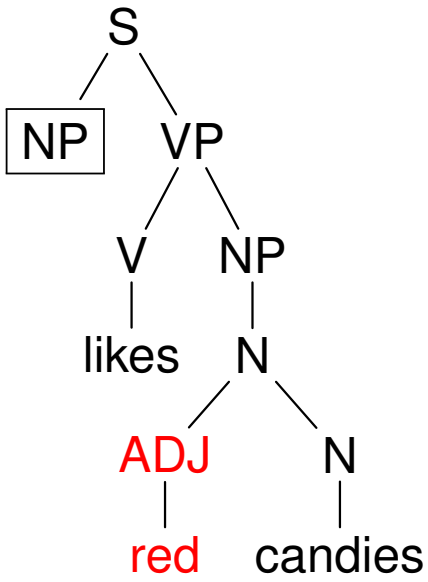

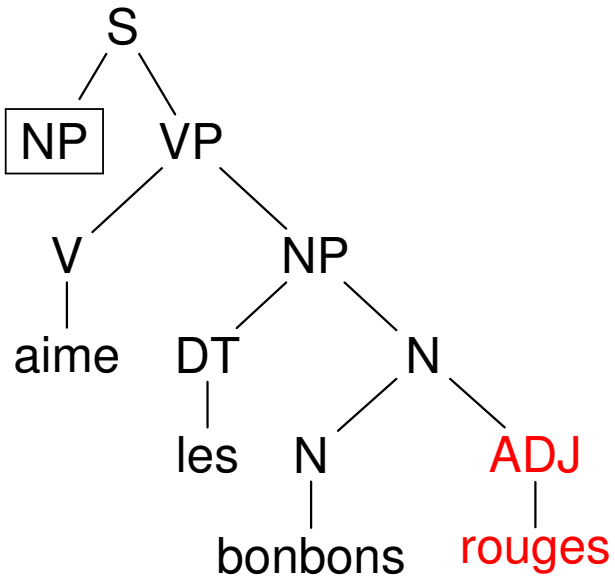

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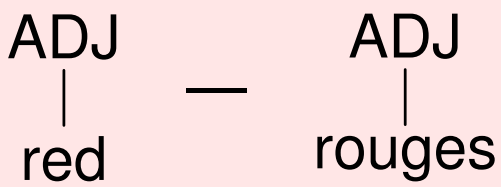


Main Question

Theorem

Every STSG is an STAG.

Question

Are they further related?

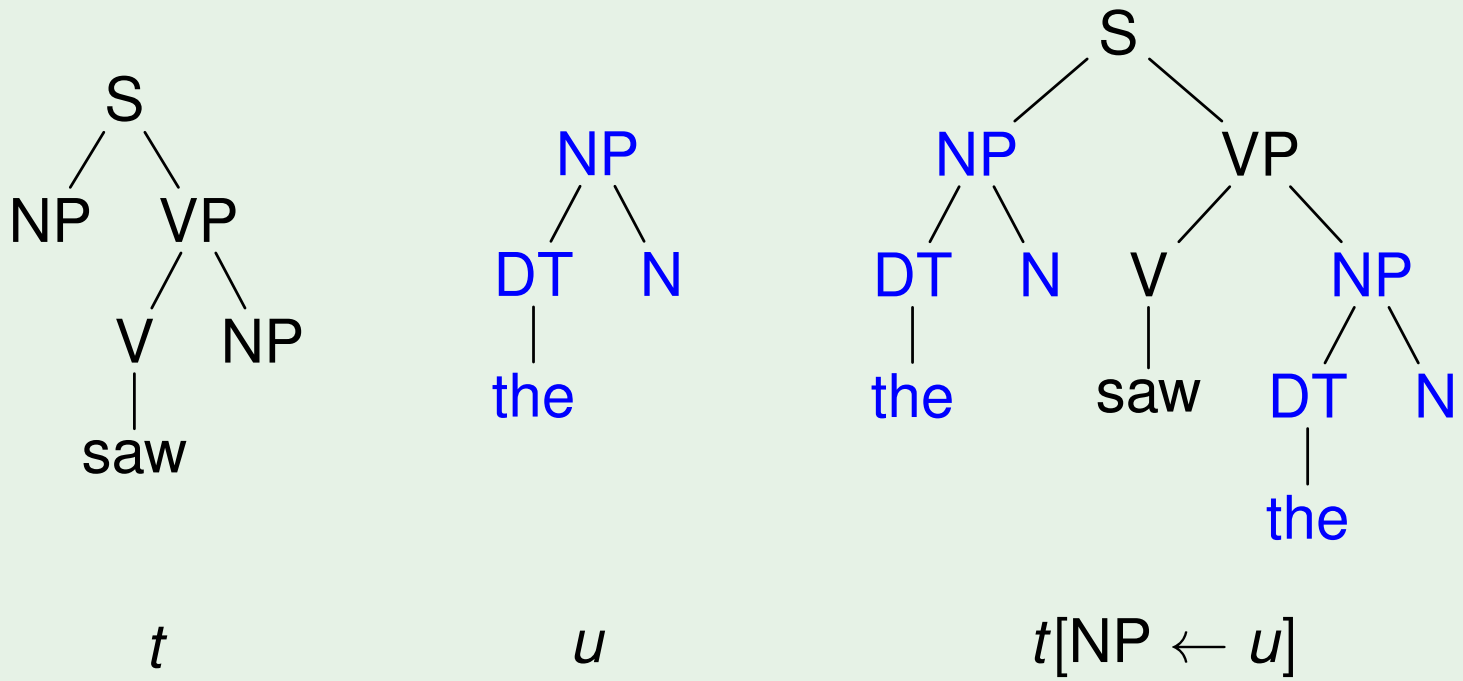
Roadmap

First-Order Substitution

Definition

$t[v_1 \leftarrow t_1, \ldots, v_k \leftarrow t_k]$ denotes the result obtained by replacing (in parallel) all occurrences of leaves labelled v_i in t by t_i.

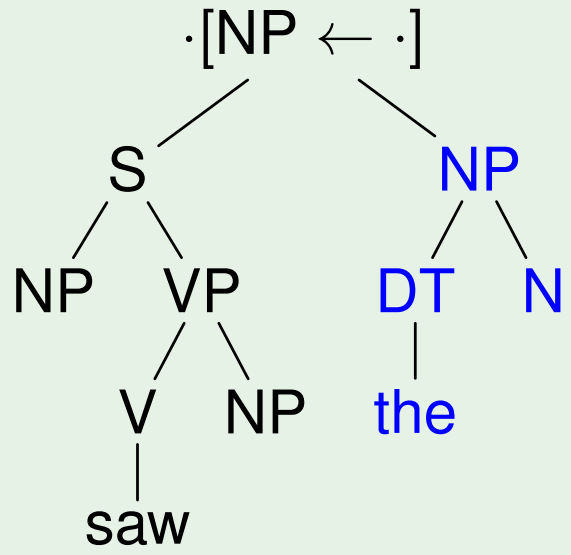
Example



t u $t[\text{NP} \leftarrow u]$

Second-Order Substitution

Example

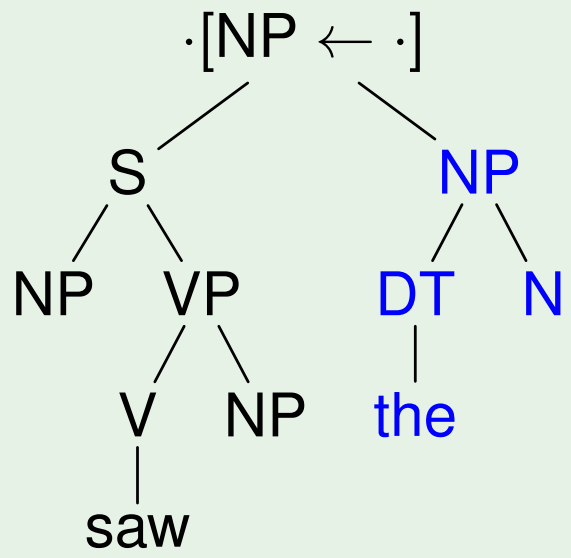


Explicit substitution

- keep an explicit representation of substitutions in tree
- any number of substitutions allowed at any level

Second-Order Substitution

Example



Evaluation

$$\text{eval}(\cdot[x \leftarrow \cdot](t, u)) = \text{eval}(t)[x \leftarrow \text{eval}(u)]$$

$$\text{eval}(\sigma(t_1, \ldots, t_k)) = \sigma(\text{eval}(t_1), \ldots, \text{eval}(t_k))$$

Second-Order Substitution

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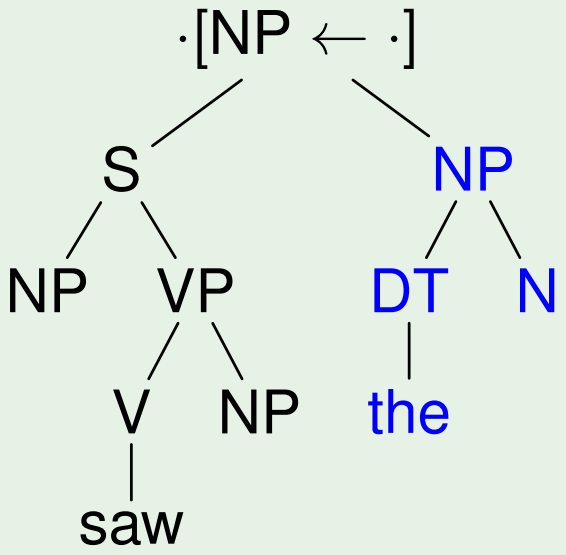

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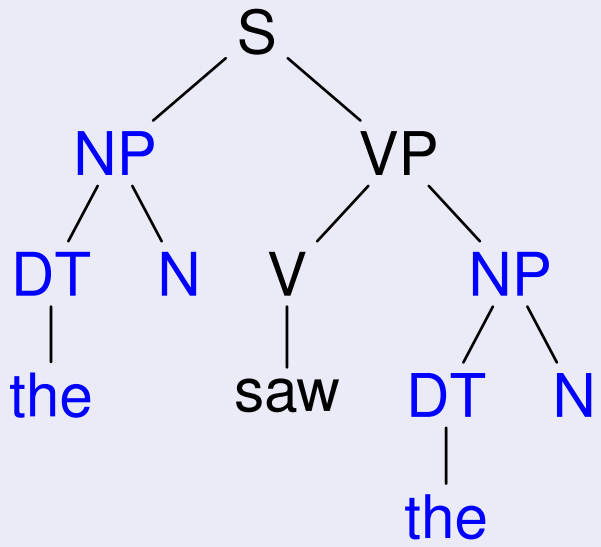

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Roadmap

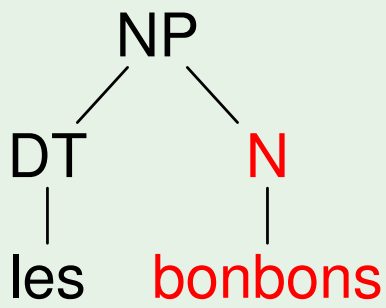
Tree-Adjoining Grammar

Intuition

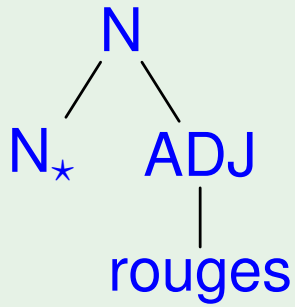
A TAG has two types of rules:

- substitution rules (as in TSG)
- adjunction rules

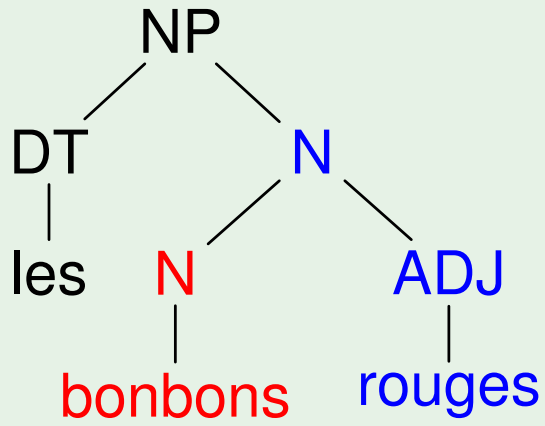
Example (Adjunction)



derived tree



auxiliary tree



adjunction

Tree-Adjoining Grammar (cont'd)

Simplifications (see [SHIEBER '06])

- no substitution rules
- adjunction mandatory (if possible)
- each adjunction spot used at most once
- root nodes of auxiliary trees are never adjunction spots

Definition

A TAG is a finite set of

- derived trees (initial trees) and
- auxiliary trees (those containing a starred node)

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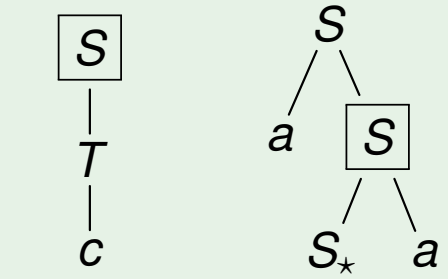
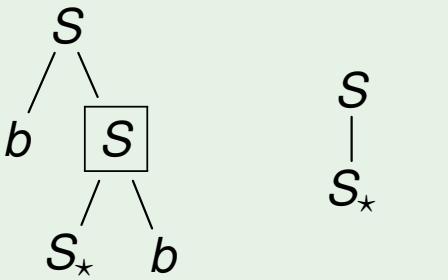


initial tree

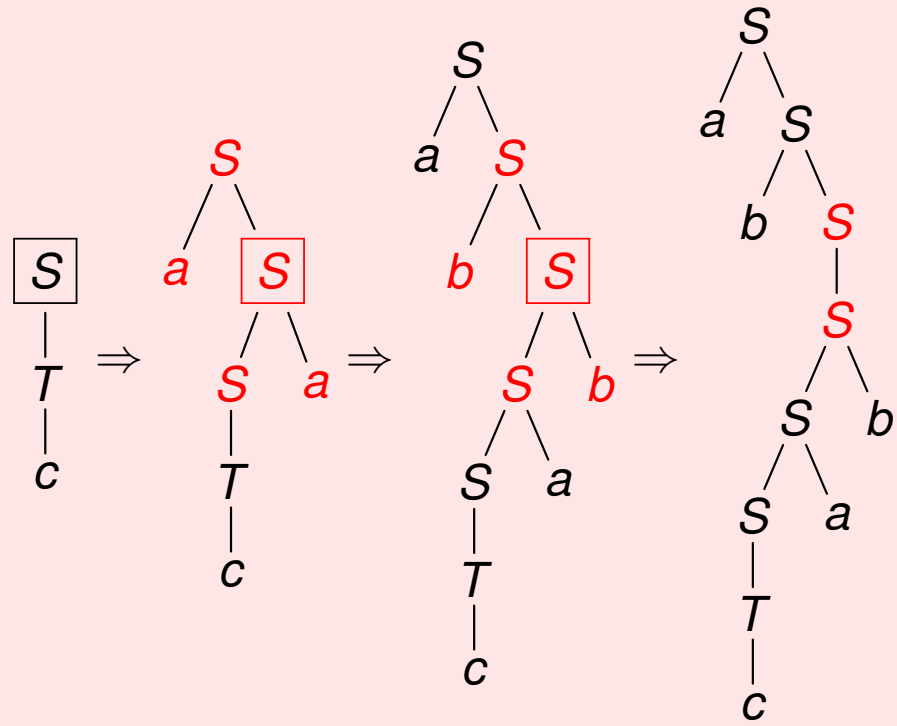
auxiliary tree



auxiliary tree

auxiliary tree

Derivation

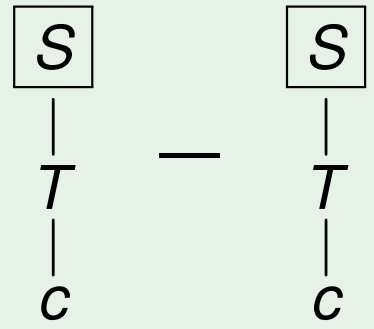


String language

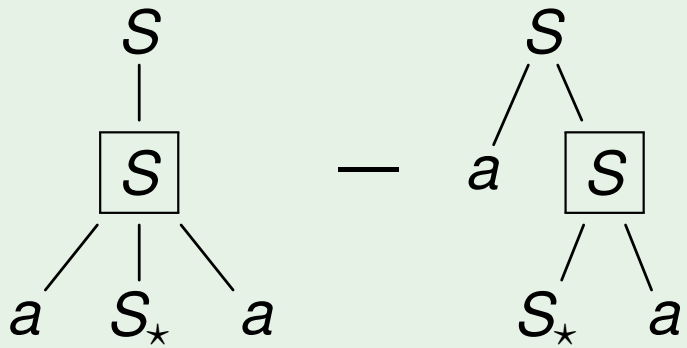
$$\{wcw \mid w \in \{a, b\}^*\}$$

Synchronous Tree-Adjoining Grammar

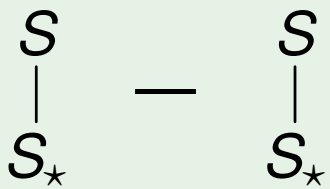
Example



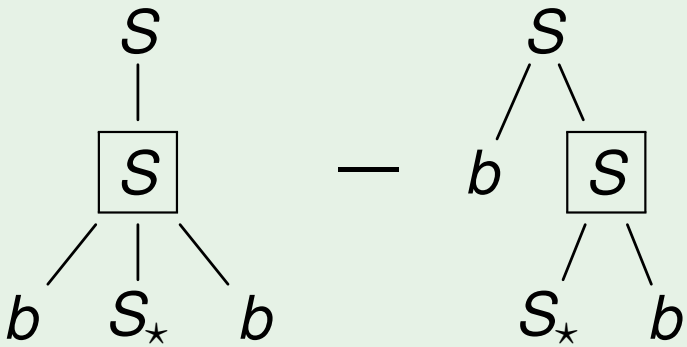
initial tree pair



auxiliary tree pair



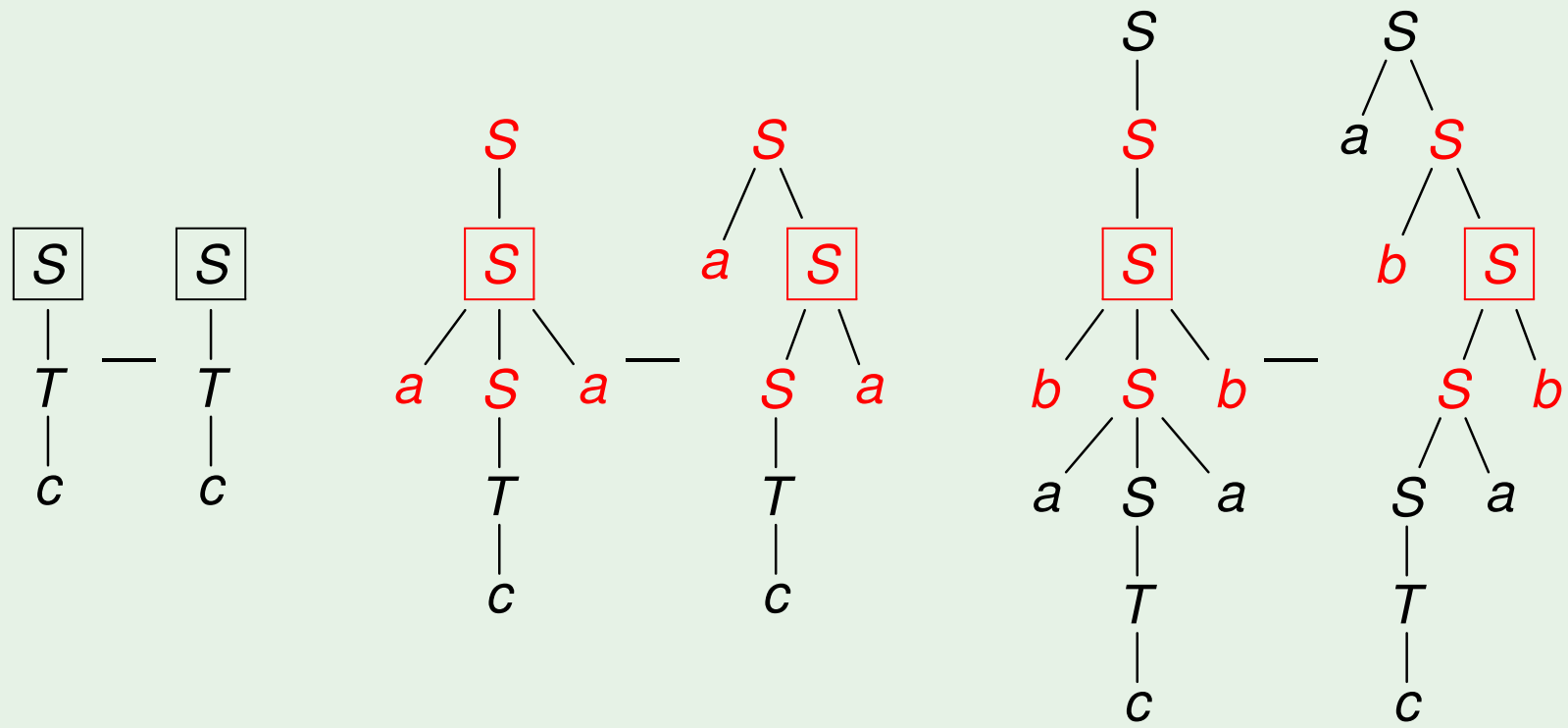
auxiliary tree pair



auxiliary tree pair

Synchronous Tree-Adjoining Grammar (cont'd)

Example



String translation

$$\{(wcw^{\mathrm{R}}, wcw) \mid w \in \{a, b\}^*\}$$

Roadmap

Simulation

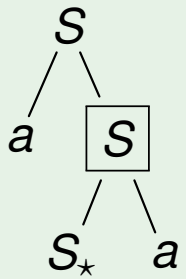
Question

Can we simulate an STAG by some STSG?

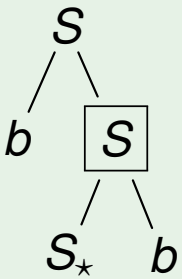
Simulation of Adjunction

Example (TAG)

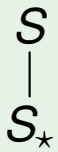
initial tree



auxiliary tree

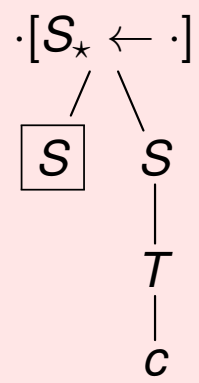


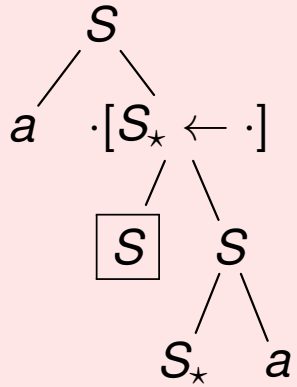
auxiliary tree

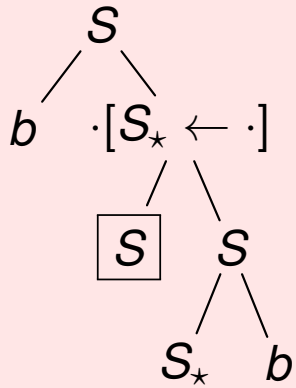


auxiliary tree

Correspondence (TSG)

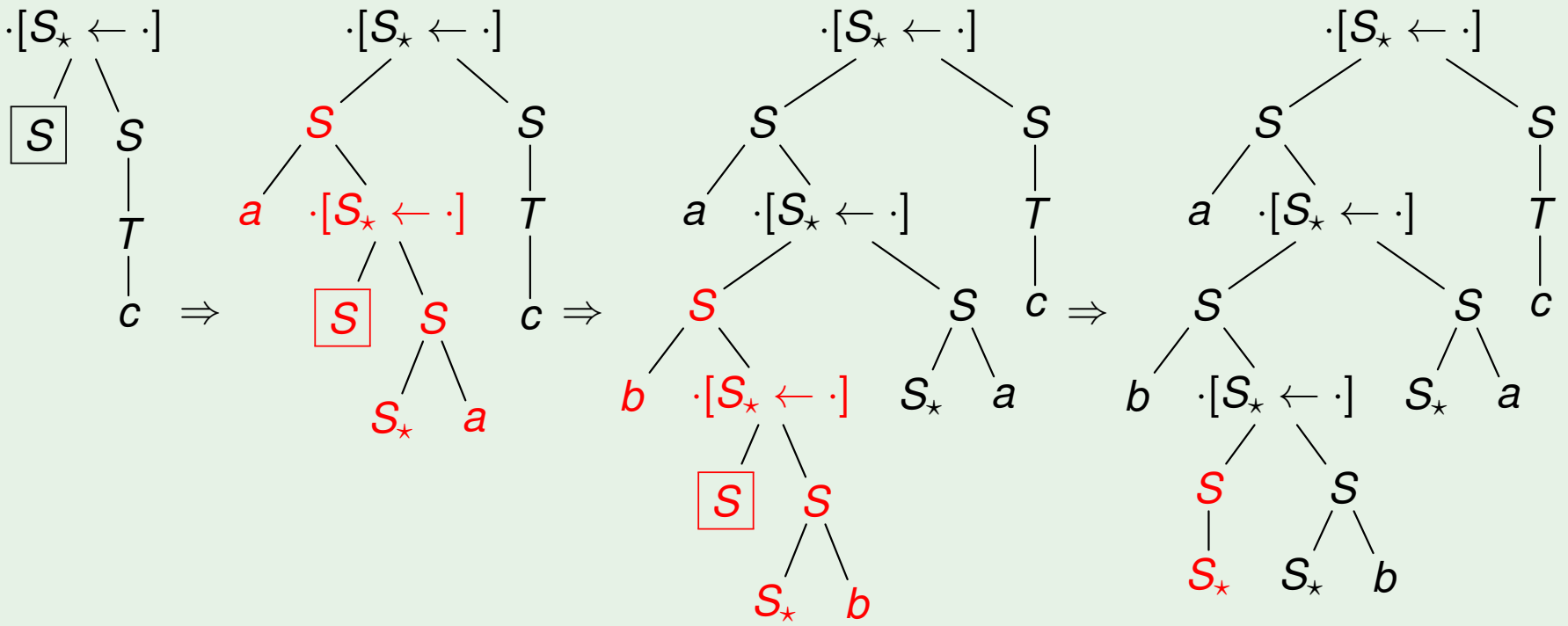






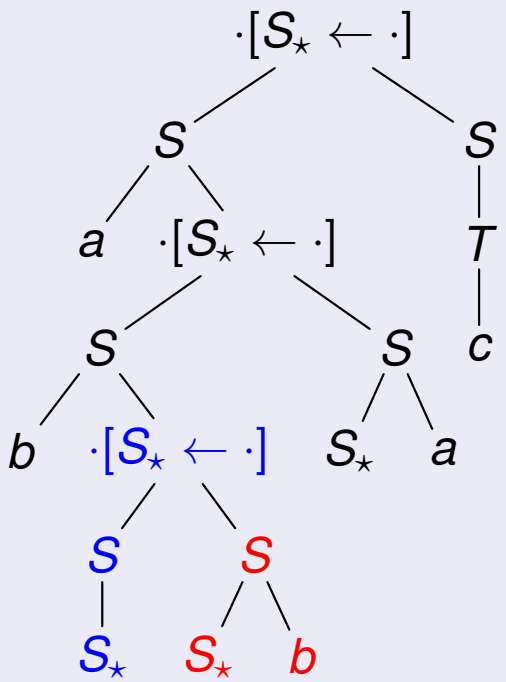
Simulation of Adjunction (cont'd)

Example

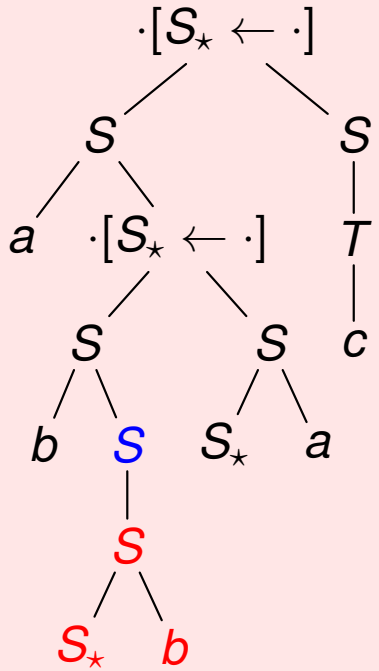


Simulation of Adjunction (cont'd)

TSG result



Evaluation

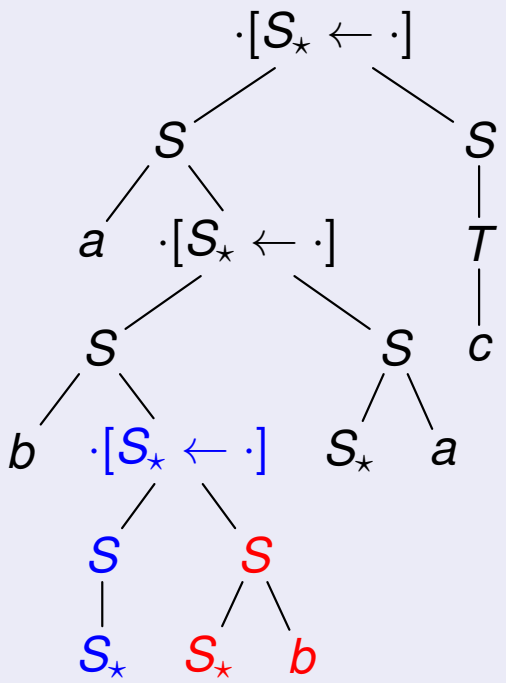


Note

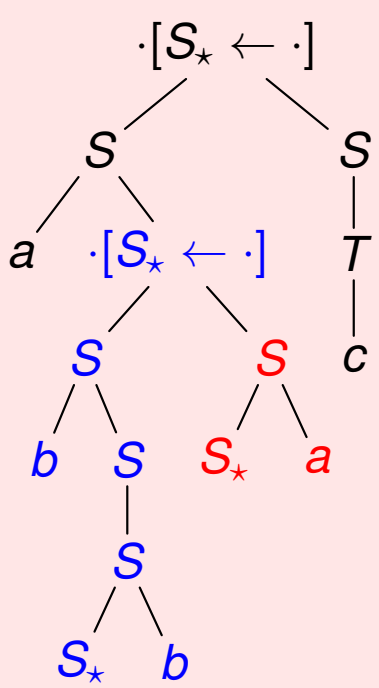
coincides with the result obtained by TAG

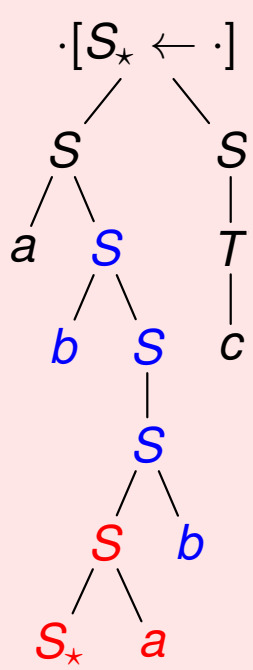
Simulation of Adjunction (cont'd)

TSG result



Evaluation





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Simulation of Adjunction (cont'd)

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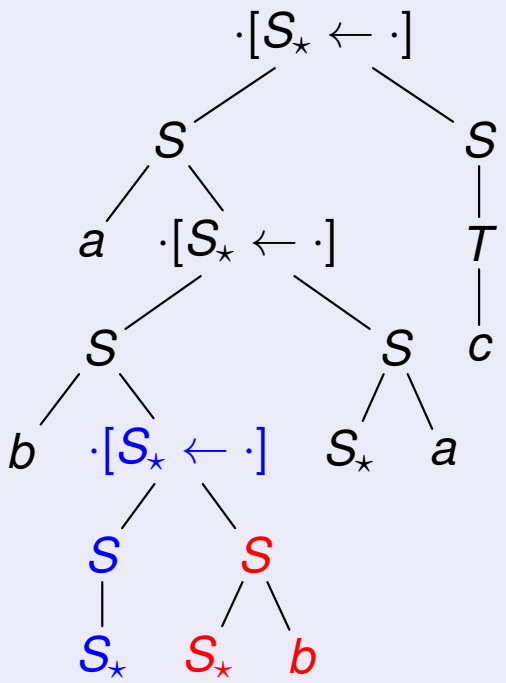

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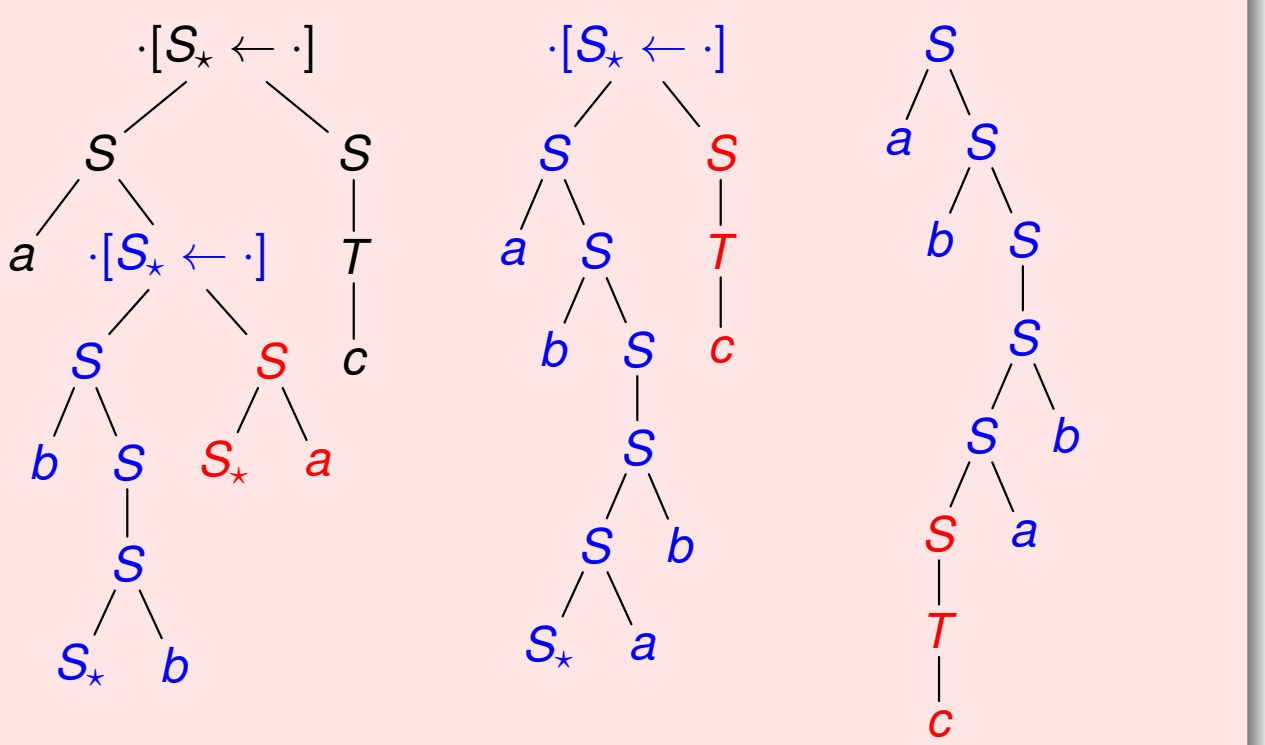

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Main Result

Theorem

For every TAG G there exists a TSG G′ such that

$$L(G) = \{\text{eval}(t) \mid t \in L(G')\}$$

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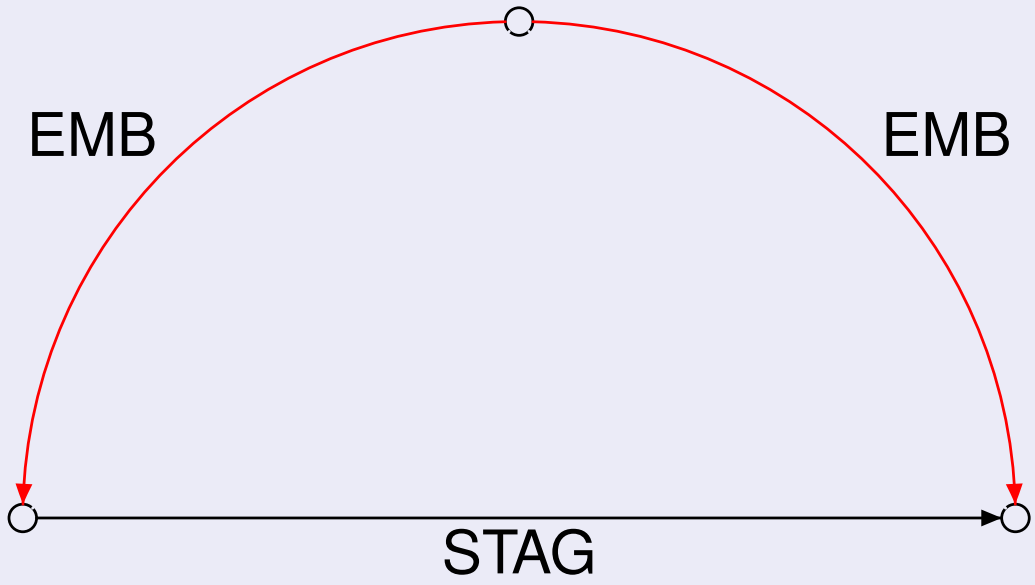

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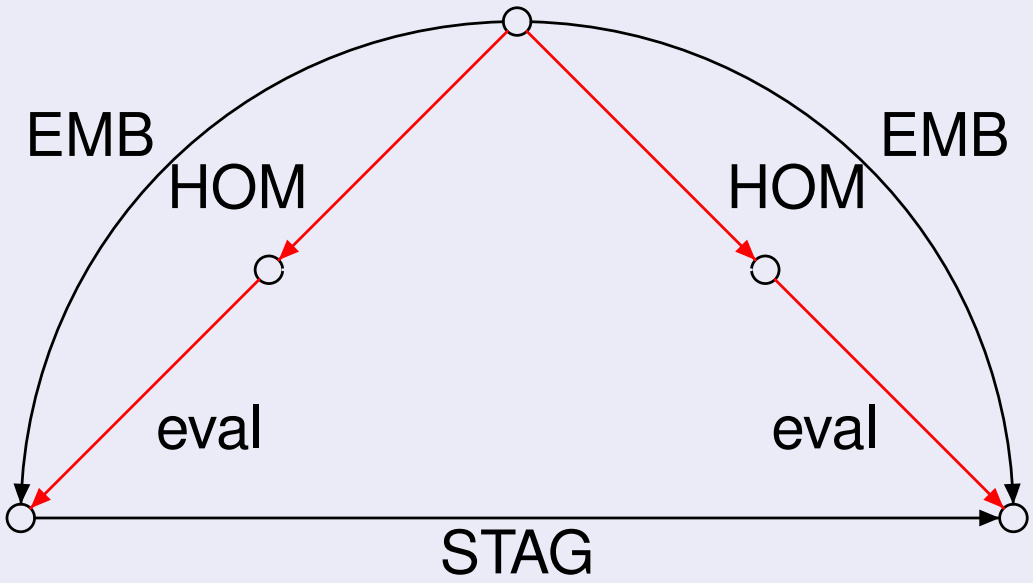

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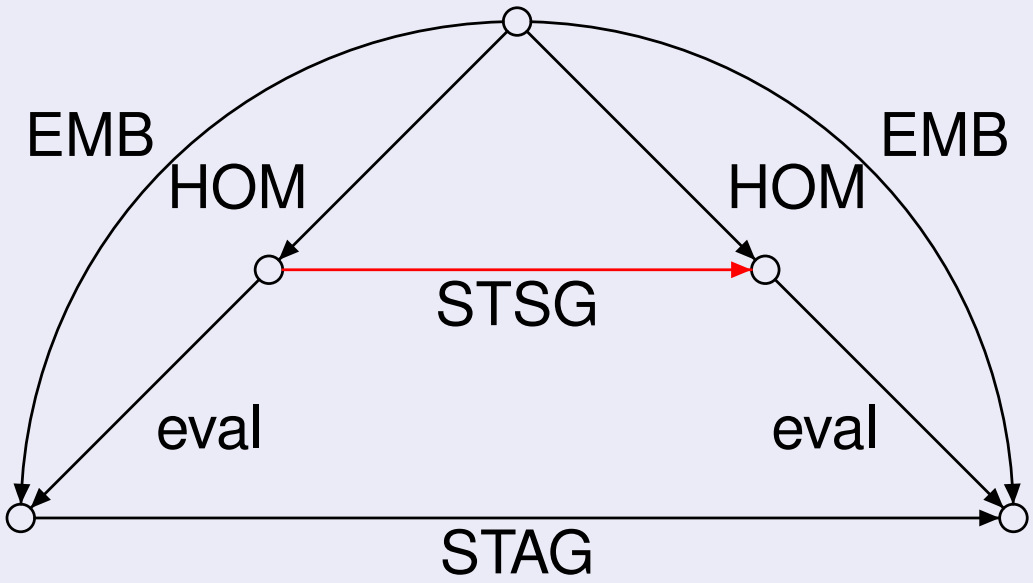

Roadmap

Application

Overview

- run an STAG in TIBURON (which can run STSGs)
- translate STSG algorithms to STAGs (factorization, etc.)
- integrate explicit substitution into semantics
- separate “context-free” and “context-sensitive” behavior

References

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- SHIEBER, SCHABES: Synchronous tree-adjoining grammars. *Computational Linguistics* 3. 1990
- SHIEBER: Unifying synchronous tree adjoining grammars and tree transducers via bimorphisms. In *EACL*. 2006

Thank you for your attention!