

Hierarchies of Tree Series Transducers—Revisited¹

Andreas Maletti²

Technische Universität Dresden
Fakultät Informatik

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1 Machine Translation

2 Tree Series Transducers

3 Results

Machine Translation

Problem

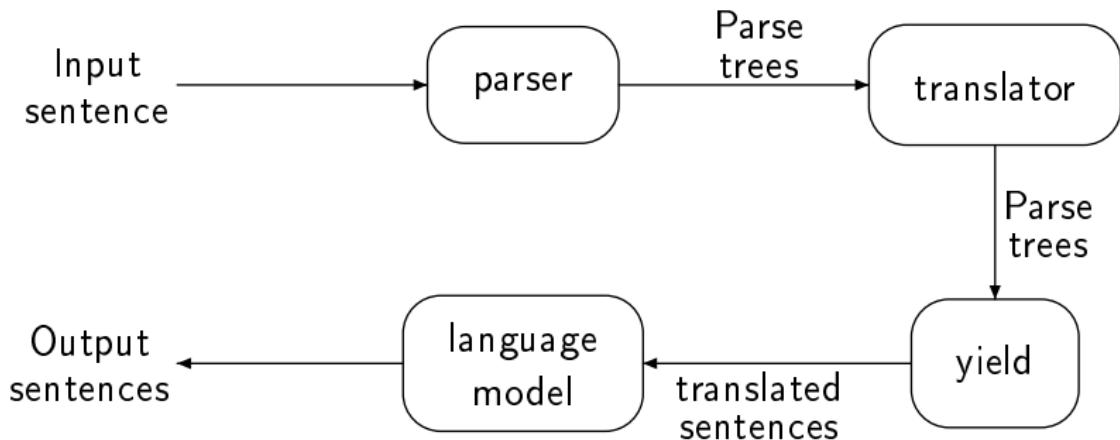
Translate text of language X into grammatical text of language Y.

- ① Preserve meaning
- ② Preserve connotation
- ③ Preserve style

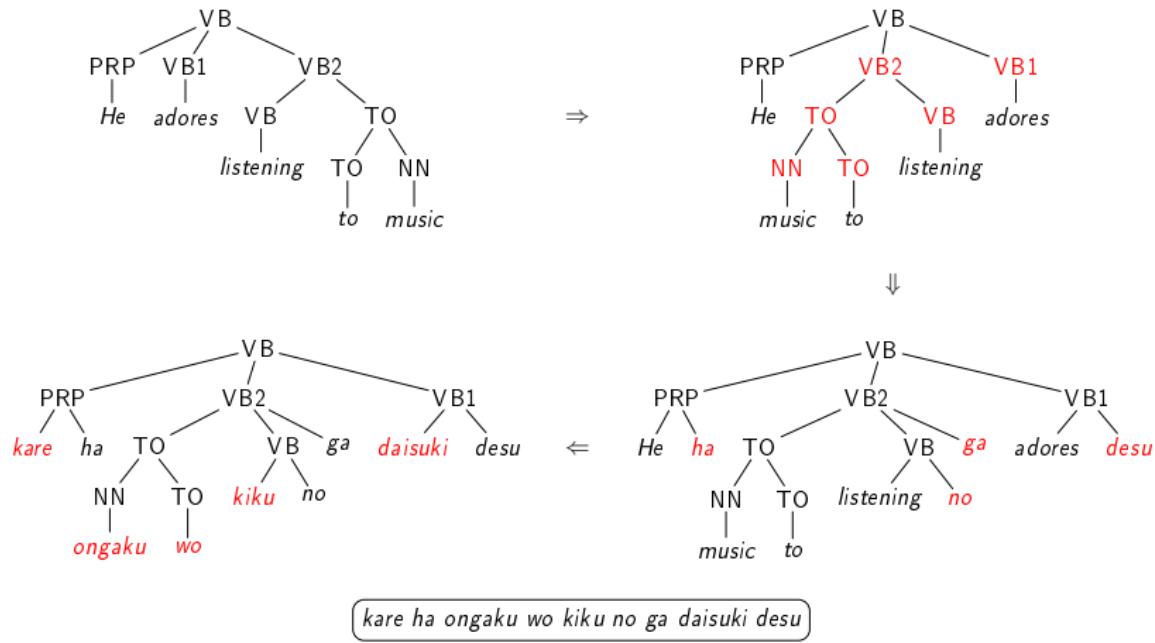
Approaches

- **traditional** phrase-based systems (e.g. CARMEL)
- **upcoming** syntax-based systems (e.g. Tiburon)

Syntax-based Machine Translation System



Syntax-based Model [Yamada, Knight 01]



3 phases: (i) Reorder, (ii) Insert, (iii) Translate

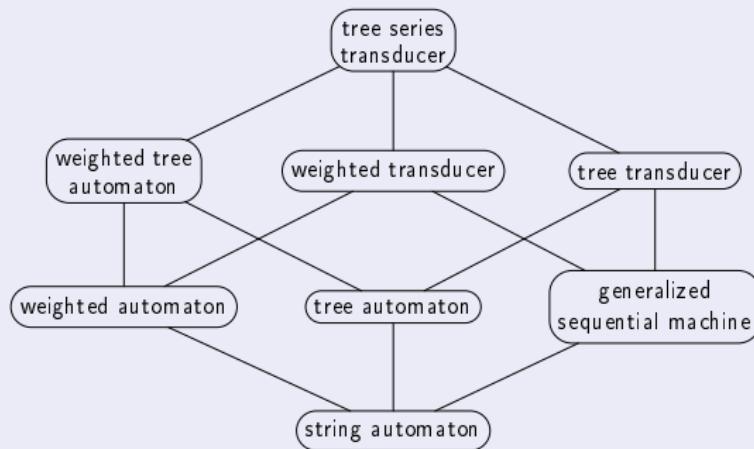
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Tree Series Transducers

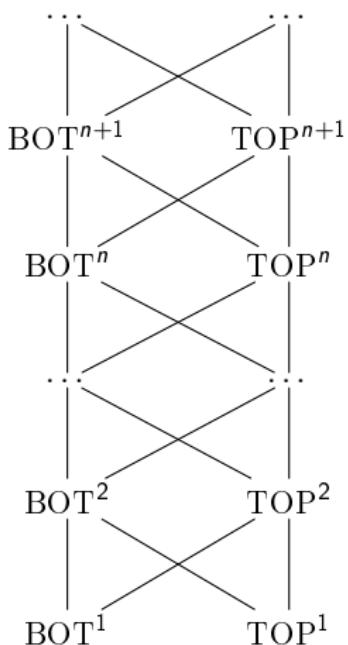
Overview



History

- Introduced in [Kuich 99]
- Extended to full generality in [Engelfriet, Fülöp, Vogler 02]

Hasse Diagram for Compositions [Engelfriet 82]



Semiring

Definition

$(A, +, \cdot, 0, 1)$ **semiring**, if

- $(A, +, 0)$ commutative monoid
- $(A, \cdot, 1)$ monoid
- \cdot distributes (both sided) over $+$
- 0 is absorbing for \cdot ($a \cdot 0 = 0 = 0 \cdot a$)

Example

- Reals $(\mathbb{R}, +, \cdot, 0, 1)$
- Probabilities $([0, 1], \max, \cdot, 0, 1)$
- Subsets $(\mathcal{P}(A), \cup, \cap, \emptyset, A)$

Tree Series

Definition

- Tree series is mapping $\psi: T_\Sigma(X) \rightarrow A$
- Set of tree series $A\langle\langle T_\Sigma(X)\rangle\rangle$

Conventions

- A usually endowed with semiring structure
- $\widetilde{0}$ is tree series that maps every tree to 0
- $\psi(t)$ written as (ψ, t)

Tree Series

Definition

Let $\psi \in A\langle\langle T_\Sigma(X) \rangle\rangle$

$$\text{supp}(\psi) = \{t \in T_\Sigma(X) \mid (\psi, t) \neq 0\}$$

i.e. set of nonzero-weighted trees

Convention

- tree series ψ written as $\sum_{t \in T_\Sigma(X)} (\psi, t) t$
- $(\psi + \varphi, t) = (\psi, t) + (\varphi, t)$
- $(a \cdot \psi, t) = a \cdot (\psi, t)$

Tree Series Transducer [Engelfriet et al 02]

Definition

Top-down tree series transducer $(Q, \Sigma, \Delta, \mathcal{A}, I, R)$ where

- Q finite set of states
- Σ and Δ input and output ranked alphabet
- $\mathcal{A} = (A, +, \cdot, 0, 1)$ semiring
- $I \subseteq Q$ set of initial states
- R finite set of rules of the form

$$q(\sigma(x_1, \dots, x_k)) \xrightarrow{a} t$$

where $t \in T_\Delta(Q(X_k))$

Tree Series Transducer [Engelfriet et al 02]

Definition

Bottom-up tree series transducer $(Q, \Sigma, \Delta, \mathcal{A}, F, R)$ where

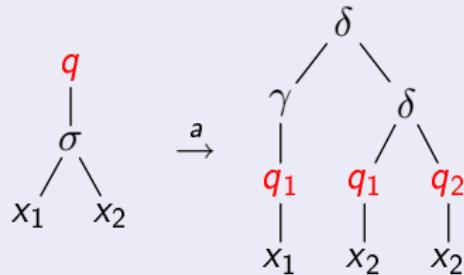
- Q finite set of states
- Σ and Δ input and output ranked alphabet
- $\mathcal{A} = (A, +, \cdot, 0, 1)$ semiring
- $F \subseteq Q$ set of final states
- R finite set of rules of the form

$$\sigma(q_1(x_1), \dots, q_k(x_k)) \xrightarrow{a} q(t)$$

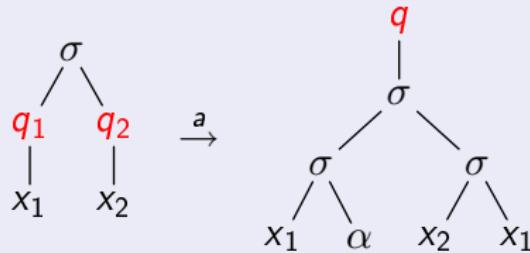
where $t \in T_\Delta(X_k)$

Tree Series Transducer

Top-down Rules

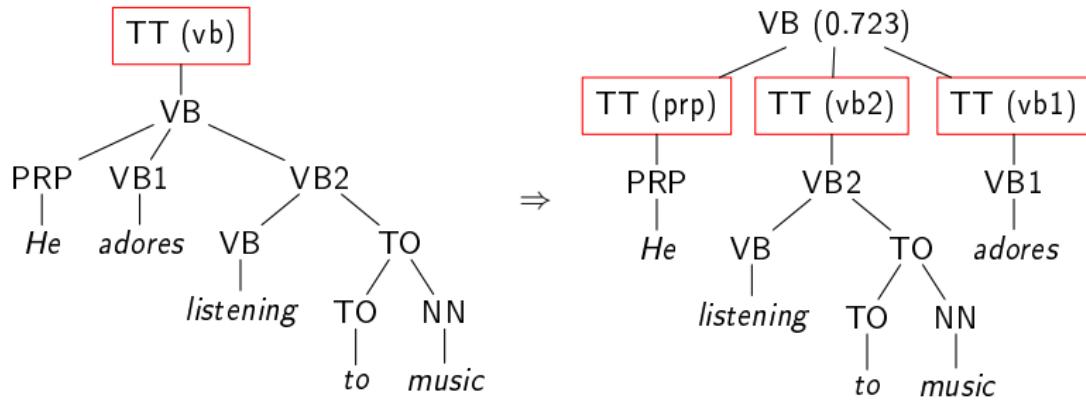


Bottom-up Rules



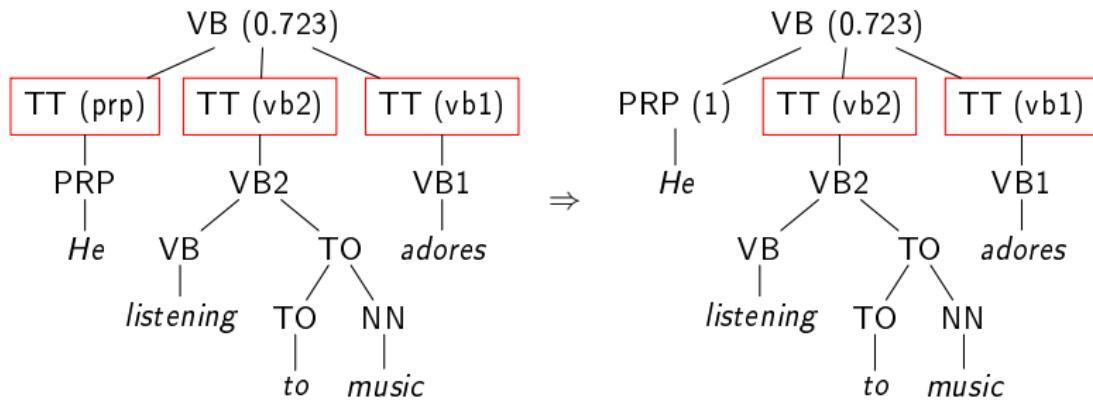
Implementation of Phase (i): Reorder

Implementation by top-down tree series transducer



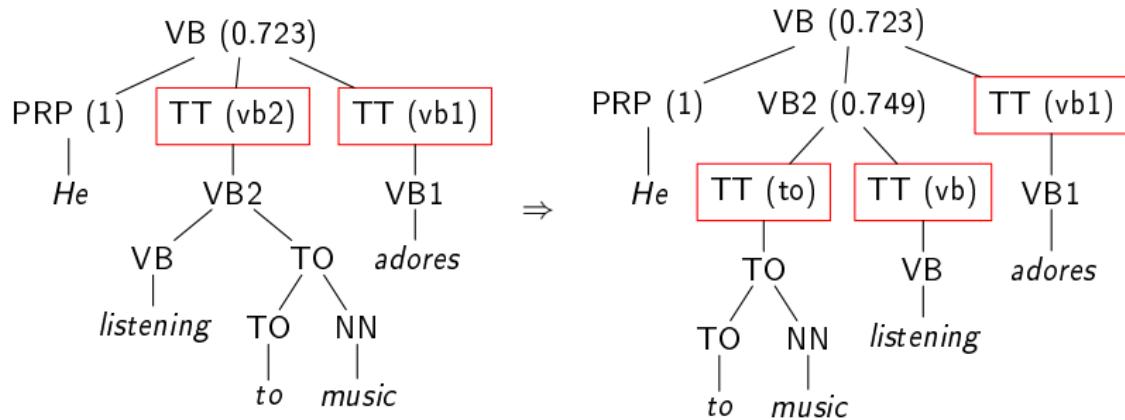
$$\text{vb}(\text{VB}(x_1, x_2, x_3)) \xrightarrow{0.723} \text{VB}(\text{prp}(x_1), \text{vb2}(x_3), \text{vb1}(x_2))$$

Implementation of Phase (i): Reorder



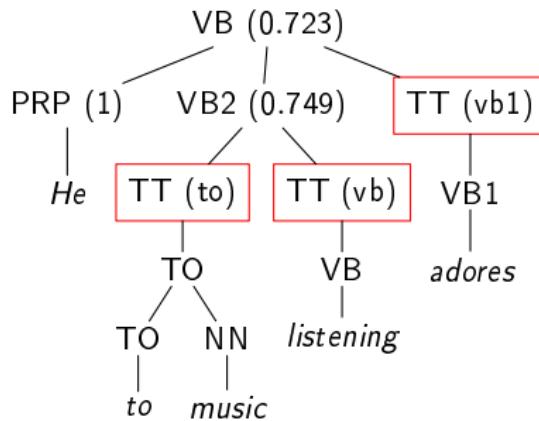
$$\text{prp}(\text{PRP}(x_1)) \xrightarrow{1} \text{PRP}(x_1)$$

Implementation of Phase (i): Reorder

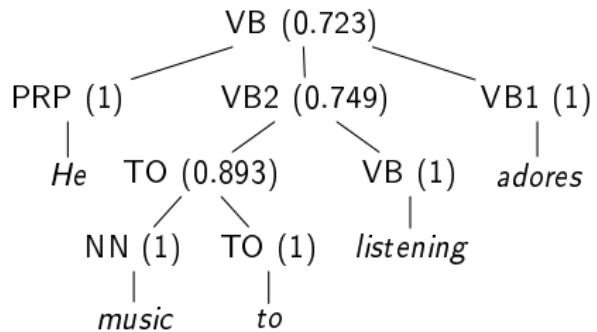


$$\text{vb2}(\text{VB2}(x_1, x_2)) \xrightarrow{0.749} \text{VB2}(\text{to}(x_2), \text{vb}(x_1))$$

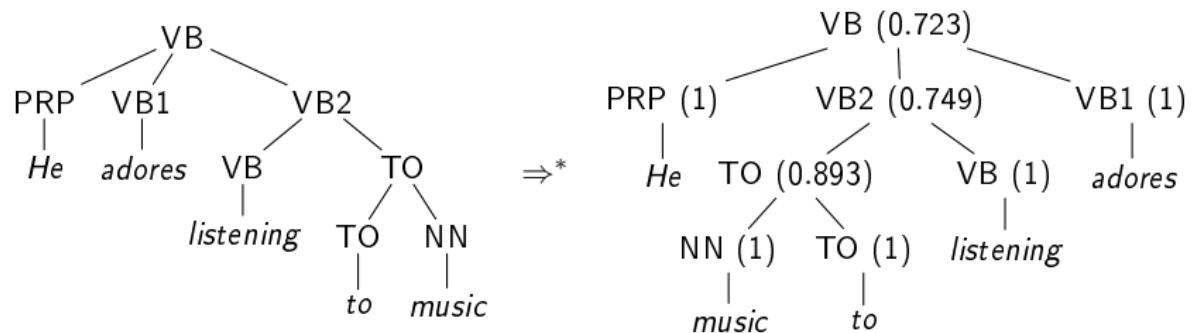
Implementation of Phase (i): Reorder



\Rightarrow^*



Implementation of Phase (i): Reorder



The above reordering has probability:

$$0.723 \cdot 0.749 \cdot 0.893 = 0.484$$

Classes of Transformations

Definition

denotation	class of transformations computed by
$\text{TOP}(\mathcal{A})$	top-down tree series transducers
$\text{BOT}(\mathcal{A})$	bottom-up tree series transducers

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Composition of Transformations

Definition

Let

- $\varphi: T_\Sigma \times T_\Delta \rightarrow A$
- $\psi: T_\Delta \times T_\Gamma \rightarrow A$

Composition of φ and ψ

$$(\varphi; \psi): T_\Sigma \times T_\Gamma \rightarrow A$$

$$(t, v) \mapsto \sum_{u \in T_\Delta} \varphi(t, u) \cdot \psi(u, v)$$

Required Semiring Properties

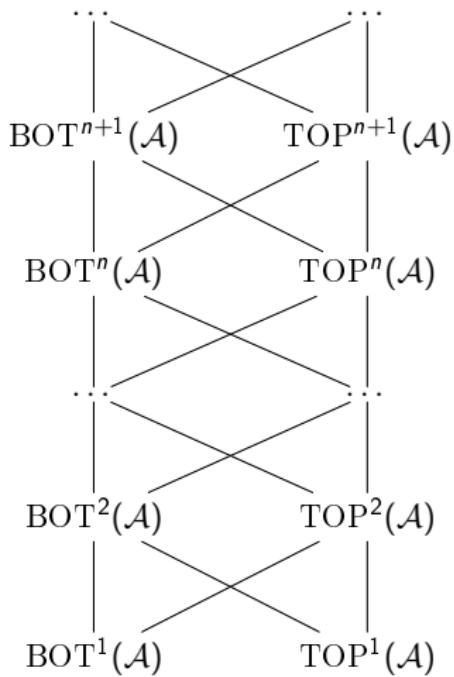
Definition

$$\mathcal{A} = (A, +, \cdot, 0, 1)$$

- **commutative**, if $a \cdot b = b \cdot a$
- **idempotent**, if $1 + 1 = 1$
- **zero-divisor free**, if $a \cdot b = 0$ implies $a = 0$ or $b = 0$
- **zero-sum free**, if $a + b = 0$ implies $a = 0$
- **positive**, if zero-divisor free and zero-sum free

Hasse Diagram for Compositions [Fülöp et al 05]

\mathcal{A} commutative, positive **idempotent** semiring



Homomorphic Images

Definition

$$\chi: A \rightarrow \{0, 1\}$$

$$\chi(a) = \begin{cases} 0 & \text{if } a = 0 \\ 1 & \text{otherwise} \end{cases}$$

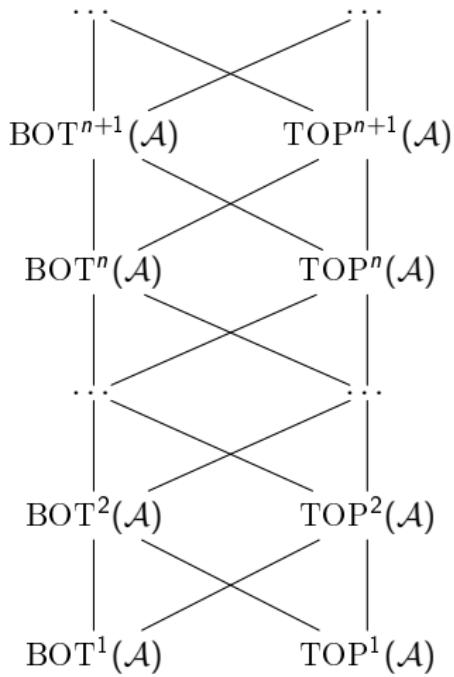
Theorem

\mathcal{A} positive semiring

- $\chi(\text{TOP}(\mathcal{A})) = \text{TOP}(\mathbb{B})$ and $\chi(\text{BOT}(\mathcal{A})) = \text{BOT}(\mathbb{B})$
- $\chi(\tau_1 ; \tau_2) = \chi(\tau_1) ; \chi(\tau_2)$

Hasse Diagram for Compositions

\mathcal{A} commutative, positive semiring (e.g. \mathbb{R}_+)



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