# Geometric Analogue of Holographic Reduced Representations 

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- Background
- Distributed representations
- Previous architectures
- GA model
- Binary parametrization of the geometric product
- Cartan representation
- Example
- Signatures
- Test results
- Recognition tests
- Blade linearity
- Future work
- Computer program CartanGA

Distributed representations

## Example

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+ "triangular shape"

$=$ "green triangle"

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## Idea

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- the size of a distributed representation is usually fixed
- the units have either binary or a continuous-space values
- data patterns are chosen as random vectors or matrices
- in most distributed representations only the overall pattern of activated units has a meaning (resemblance to a multi-superimposed photograph)

Distributed representations

Atomic objects and complex statements

Distributed representations

## Atomic objects and complex statements

```
bite agent }\circledast\mathrm{ Fido + bite object }\circledast\mathrm{ Pat = "Fido bit Pat"
    chunk chunk
- roles: bite \(_{\text {agent }}\), bite \(_{\text {object }}\)
- fillers: Fido, Pat
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name \(\circledast\) Pat + sex \(\circledast\) male + age \(\circledast 66=\) "PSmith"
    - roles: name, sex, age
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## Binding and chunking

$\circledast$ binding

+ superposition (also called chunking)


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Decoding

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$\#$ decoding
$x^{-1}$ (approximate/pseudo) inverse

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## Clean-up memory

- An auto-associative collection of elements and statements (excluding single chunks) produced by the system.
- Given a noisy extracted vector such structure must be able to:
- either recall the most similar item stored
- ... or indicate, that no matching object had been found.
- We need a measure of similarity (in most models: scalar product, Hamming/Euclidean distance).


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## Advantages

- Noise tolerance, good error correction properties.
- Storage efficiency.
- The number of superimposed patterns is not fixed.


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- Similarity measure: dot product

Binary parametrization of the geometric product

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Let $\left\{x_{1}, \ldots, x_{n}\right\}$ be a string of bits and $\left\{e_{1}, \ldots, e_{n}\right\}$ be base vectors. Then

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e_{x_{1} \ldots x_{n}}=e_{1}^{x_{1}} \ldots e_{n}^{x_{n}}
$$

where $e_{k}^{0}=1$.

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## Geometric product as a projective XOR representation

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$$
\begin{aligned}
e_{12} e_{1} & =e_{110 \ldots 0} e_{10 \ldots 0}=e_{1} e_{2} e_{1}=-e_{2} e_{1} e_{1}=-e_{2}=-e_{010 \ldots 0} \\
& =-e_{(110 \ldots 0) \oplus(10 \ldots 0)}=(-1)^{D} e_{(110 \ldots 0) \oplus(10 \ldots 0)}
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More formally, for two arbitrary strings of bits we have:

$$
e_{A_{1} \ldots A_{n}} e_{B_{1} \ldots B_{n}}=(-1)^{\sum_{k<1} B_{k} A_{l}} e_{\left(A_{1} \ldots A_{n}\right) \oplus\left(B_{1} \ldots B_{n}\right)}
$$

## Matrix representation

Pauli's matrices

$$
\sigma_{1}=\left(\begin{array}{ll}
0 & 1 \\
1 & 0
\end{array}\right), \quad \sigma_{2}=\left(\begin{array}{cc}
0 & -i \\
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## Cartan representation

$$
\begin{aligned}
b_{2 k} & =\underbrace{\sigma_{1} \otimes \cdots \otimes \sigma_{1}}_{n-k} \otimes \sigma_{2} \otimes \underbrace{1 \otimes \cdots \otimes 1}_{k-1} \\
b_{2 k-1} & =\underbrace{\sigma_{1} \otimes \cdots \otimes \sigma_{1}}_{n-k} \otimes \sigma_{3} \otimes \underbrace{1 \otimes \cdots \otimes 1}_{k-1} .
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Representation

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\text { male } & =c_{00111}=b_{3} b_{4} b_{5} \\
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& =1 \otimes 1 \otimes 1 \otimes 1 \otimes\left(-i \sigma_{1}\right) \\
\text { name } & =c_{00010}=b_{4}=\sigma_{1} \otimes \sigma_{1} \otimes \sigma_{1} \otimes \sigma_{2} \otimes 1 \\
\text { sex } & =c_{11100}=b_{1} b_{2} b_{3} \\
& =\left(\sigma_{1} \otimes \sigma_{1} \otimes \sigma_{1} \otimes \sigma_{1} \otimes \sigma_{3}\right)\left(\sigma_{1} \otimes \sigma_{1} \otimes \sigma_{1} \otimes \sigma_{1} \otimes \sigma_{2}\right)\left(\sigma_{1} \otimes \sigma_{1} \otimes \sigma_{1} \otimes \sigma_{3} \otimes 1\right) \\
& =\sigma_{1} \otimes \sigma_{1} \otimes \sigma_{3} \otimes 1 \otimes\left(-i \sigma_{1}\right) \\
\text { age } & =c_{10001}=b_{1} b_{5}=\left(\sigma_{1} \otimes \sigma_{1} \otimes \sigma_{1} \otimes \sigma_{1} \otimes \sigma_{3}\right)\left(\sigma_{1} \otimes \sigma_{1} \otimes \sigma_{3} \otimes 1 \otimes 1\right) \\
& =1 \otimes 1 \otimes\left(-i \sigma_{2}\right) \otimes \sigma_{1} \otimes \sigma_{3}
\end{aligned}
$$

## Example

Encoding and Decoding

## PSmith

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Encoding and Decoding

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## $P S m i t h=$ name $\circledast$ Pat + sex $\circledast$ male + age $\circledast 66$

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\begin{aligned}
\text { PSmith } & =\text { name } \circledast \text { Pat }+ \text { sex } \circledast \text { male }+ \text { age } \circledast 66 \\
& =c_{00010} c_{00100}+c_{11100} c_{00111}+c_{10001} c_{11000}
\end{aligned}
$$

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\begin{aligned}
\text { PSmith } & =\text { name } \circledast \text { Pat }+ \text { sex } \circledast \text { male }+ \text { age } \circledast 66 \\
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& =-c_{00100}-c_{11001}-c_{01011}
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## PSmith

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\begin{aligned}
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& =c_{00010} c_{00100}+c_{11100} c_{00111}+c_{10001} c_{11000} \\
& =-c_{00110}+c_{11011}+c_{01001}
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$$

## PSmith's name $=P$ Smith $\sharp$ name ${ }^{-1}=$ PSmith $\circledast$ name

$$
\begin{aligned}
\text { PSmith } \circledast \text { name } & =\left(-c_{00110}+c_{11011}+c_{01001}\right) c_{00010} \\
& =-c_{00100}-c_{11001}-c_{01011} \\
& =- \text { Pat }+ \text { noise }=\text { Pat }^{\prime}
\end{aligned}
$$

## Example

Clean-up

## Scalar product

Scalar (inner) product is performed by the means of matrix trace.

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$$
\begin{aligned}
\left\langle P a t^{\prime} \mid e_{P a t}\right\rangle & =\operatorname{Tr}\left(\left(-c_{00100}-c_{11001}-c_{01011}\right) c_{00100}\right) \\
& =\operatorname{Tr}\left(-1+c_{11101}-c_{01111}\right) \\
& =-32 \neq 0
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& =\operatorname{Tr}\left(-\mathbf{1}+c_{11101}-c_{01111}\right) \\
& =-32 \neq 0 \\
\left\langle P a t^{\prime} \mid e_{\text {male }}\right\rangle & =0 \\
\left\langle P a t^{\prime} \mid e_{66}\right\rangle & =0 \\
\left\langle P a t^{\prime} \mid e_{\text {name }}\right\rangle & =0 \\
\left\langle P a t^{\prime} \mid e_{\text {sex }}\right\rangle & =0 \\
\left\langle P a t^{\prime} \mid e_{a g e}\right\rangle & =0 \\
\left\langle P a t^{\prime} \mid P S m i t h\right\rangle & =0
\end{aligned}
$$

## Signatures

Examples

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Examples


## Signatures

Examples


## Signatures

Dimensions

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- $2^{\left\lfloor\frac{n}{2}\right\rfloor}$ signatures on one of the diagonals.
- Each signature is of dimensions $2^{\left\lceil\frac{n}{2}\right\rceil} \times 2^{\left\lceil\frac{n}{2}\right\rceil}$.


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## Cartan representation

$$
\begin{aligned}
b_{2 k} & =\underbrace{\sigma_{1} \otimes \cdots \otimes \sigma_{1}}_{\text {at least }\left\lfloor\frac{n}{2}\right\rfloor} \otimes \sigma_{2} \otimes \underbrace{1 \otimes \cdots \otimes 1}_{k-1}, \\
b_{2 k-1} & =\underbrace{\sigma_{1} \otimes \cdots \otimes \sigma_{1}}_{\text {at least }\left\lfloor\frac{n}{2}\right\rfloor} \otimes \sigma_{3} \otimes \underbrace{1 \otimes \cdots \otimes 1}_{k-1} .
\end{aligned}
$$

## Recognition tests

## Construction

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- Plate (HRR) construction, e.g. eat + eat $_{\text {agent }} \circledast$ Mark + eat $_{\text {object }} \circledast$ theFish


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- Plate (HRR) construction, e.g. eat + eat agent $\circledast$ Mark + eat $t_{\text {object }} \circledast$ theFish
- Agent-Object construction, e.g. eat $t_{\text {agent }} \circledast$ Mark + eat object $\circledast$ theFish


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- Plate (HRR) construction, e.g. eat + eat agent $\circledast$ Mark + eat $t_{\text {object }} \circledast$ theFish
- Agent-Object construction, e.g. eat $t_{\text {agent }} \circledast$ Mark + eat object $\circledast$ theFish


## Vocabulary/Sentence set

Similar sentences, e.g:

- Fido bit Pat.
- Fido bit PSmith.
- Pat fled from Fido.
- PSmith fled from Fido.
- Fido bit PSmith causing PSmith to flee from Fido.
- Fido bit Pat causing Pat to flee from Fido.
- John saw that Fido bit PSmith causing PSmith to flee from Fido.
- John saw that Fido bit Pat causing Pat to flee from Fido.
- etc...

Altogether 42 atomic objects and 19 sentences.

## Recognition tests

## Agent-Object construction

Works better for:

- simple sentences and simple answers


## Recognition tests

## Agent-Object construction

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## Recognition tests

## Agent-Object construction

Works better for:

- simple sentences and simple answers,
- nested sentences, from which a rather unique information is to be derived.


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## Recognition tests

## Plate (HRR) construction

Works better for nested sentences, from which a complex information needs to be derived.


## Blade linearity

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## Assumption: ideal conditions

No two chunks of a sentence at any time are identical (up to a constant).

$$
\langle A \mid C\rangle \neq 0 \equiv A \text { and } C \text { share a common blade }
$$

## Blade linearity

## One meaningful blade, $L>0$ noisy blades

avg number of probings giving
a nonzero matrix trace


## Blade linearity

## $K>1$ meaningful blades, $L>0$ noisy blades

avg number of probings giving
a nonzero matrix trace*


## What's next?

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- Decide on construction (something between Plate and Agent-Object?).
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- Scaling.
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