Formal Verification Using Timed Automata

Formal Verification Using Timed Automata



The Enhanced Easy Link (EEL) Protocol

- Philips standard
- Rules for sending control information between components of Philips audio systems
- Timing is essential part of the protocol
- Philips uses one processor chip for several tasks
- Philips allows high tolerances on timing
- What is the maximal timing tolerance so that the protocol is correct?

Informal Description of EEL

- Manchester Encoding of bit strings
- Receiver must recognize first bit of a new message: bus voltage low when no message is sent, first bit of a message must be 1
- Receiver cannot recognize downgoing edges due to hardware restrictions:
 - receiver only takes into account upgoing edges
 - bit strings must be 5 + 8n bits long for some $n \in \mathbb{N}$.
- Receiver does not know length of bit string: receiver stops decoding after certain time has passed since last upgoing edge
- Receiver should not glue together two messages: waiting time before last upgoing edge and the first upgoing edge.

- Receiver keeps track of last received bit and the time elapsed since last upgoing edge.

Informal Description of EEL

- EEL is a bus protocol: all components are connected to one cable
- Uses broadcast communication
- Collision: two senders start to send at the same time
 - each sender must check whether it is the only one that is sending
 - if bus is on 1, while it sends 0, then it stops sending
 - bus level check:
 - (A) just before every upgoing edge,
 - (B) at 25% and 75% of a bit slot when sending a 0 $\,$

Informal Description of EEL

- One bit slot equals 888 microseconds
- EEL has high timing tolerance:
 - Timing of edges: $\pm 5\%$
 - Bus level check (A): \pm 20 microseconds
 - Bus level check (B): \pm 22 microseconds

- In a former version, the bus level check (A) was required only before the first edge of a message

- Bad scenario: fast sender, slow sender, perfect receiver:



- This version of the protocol was implemented and sold!

- Philips is interested in the maximal time tolerance

- In 1994, Griffioen (G94) proved the correctness of EEL for all time tolerances less than $\pm 5.35\%$ using formal methods

(G94) Griffioen: Analysis of an Audio Control Protocol with Bus Collision. 1994

Formal Methods

Writing formal specifications instead of natural language specifications
May clarify unclear sections or ambiguities, because formalisms have a formally defined semantics

- Formalisms are easy to read
- One may prove properties about the specifications, e.g., the correctness

- Philips is interested in the maximal time tolerance

- In 1994, Griffioen (G94) proved the correctness of EEL for all time tolerances less than $\pm 5.35\%$ using formal methods

- Used timed automata as a formalism
- Proved the correctness of EEL without tool support

- Philips is interested in the maximal time tolerance

- In 1994, Griffioen (G94) proved the correctness of EEL for all time tolerances less than $\pm 5.35\%$ using formal methods

- Used timed automata as a formalism
- Proved the correctness of EEL without tool support

- In 1996, Bengtsson et al (B96) *automatically* analyzed EEL using the tool *UPPAAL*

(G94) Griffioen: Analysis of an Audio Control Protocol with Bus Collision. 1994

(B96) Bengtsson et al: Verification of an Audio Protocol with Bus Collision Using UPPAAL. 1996.

Formal Methods

Writing formal specifications instead of natural language specifications
May clarify unclear sections or ambiguities, because formalisms have a formally

defined semantics

- Formalisms are easy to read
- One may prove properties about the specifications, e.g., the correctness *automatically*
- For this, the formalisms needs to have according decidable decision problems