

# THE EQUIVALENCE, UNAMBIGUITY AND SEQUENTIALITY PROBLEMS OF FINITELY AMBIGUOUS MAX-PLUS TREE AUTOMATA ARE DECIDABLE

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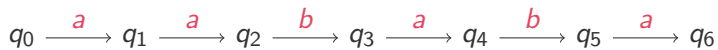
Erik Paul

Leipzig University

August 21, 2017

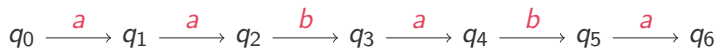


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initial weight + transition weights + final weight

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Weight of word:

maximum over all runs

# MAX-PLUS AUTOMATA: AMBIGUITY

one “initial state”

no two valid  $p \xrightarrow{a} q_1, p \xrightarrow{a} q_2$

sequential / deterministic

$$\text{Run}(w) = \{\text{Runs } r \text{ on } w \text{ with } \text{weight}(r) \neq -\infty\}$$

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$$|\text{Run}(w)| \leq 1$$



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$$|\text{Run}(w)| \leq P(|w|)$$

# THREE DECISION PROBLEMS

unambiguous	$ \text{Run}(w)  \leq 1$
finitely ambiguous	$ \text{Run}(w)  \leq M$
polynomially ambiguous	$ \text{Run}(w)  \leq P( w )$

## Equivalence problem

Given  $\mathcal{A}_1, \mathcal{A}_2$

Is  $[[\mathcal{A}_1]](w) = [[\mathcal{A}_2]](w)$  for all  $w$ ?

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Decidability for max-plus automata on words

	Equivalence	Unambiguity	Sequentiality
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Krob

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... on trees **up to now**

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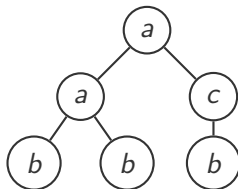
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# TREE AUTOMATA

Decidability for max-plus automata on (ranked) trees

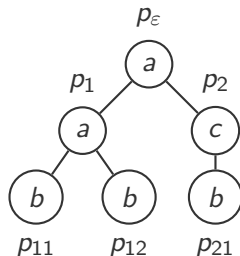
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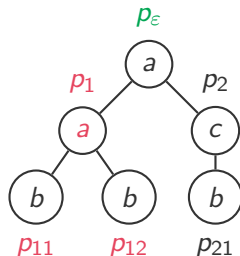
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weight of run =

transition weights + final weight

$(p_{11}, p_{12}, a, p_1)$





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$p_1$	$a$	$p_3$	$b$	$p_3$	$a$	$p_1$	$b$	$p_2$	$b$	$p_3$	$a$	$p_4$
$p_2$		$p_4$		$p_3$		$p_1$		$p_1$		$p_3$		$p_3$
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$wt_1, wt_2, wt_3 < wt_4?$

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Cycle Decomposition

$\vec{P}_1 \ x_1 \ \vec{P}_2 \ y_2 \ \vec{P}_2 \ x_3 \ \vec{P}_3 \ y_4 \ \vec{P}_3 \ x_5 \ \vec{P}_4 \ y_6 \ \vec{P}_4 \ x_7 \ \vec{P}_5$

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$x_i, y_i$  short:  $|x_i|, |y_i| \leq |\text{states}(\mathcal{A}_1)|^3 \cdot |\text{states}(\mathcal{A}_2)|$

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Vectors of weights

$$\begin{array}{rcccccccccccc} \text{wt}_1 & & 7 & & 13 & & & & 7 & & 2 & & & & 12 & & 3 & & & & 18 \\ \text{wt}_2 & = & 11 & + & 8 & & + & 3 & + & 3 & & + & 10 & + & 7 & & & + & 2 & & \\ \text{wt}_3 & & 4 & & 6 & & + & 1 & + & 15 & & + & 9 & + & 5 & & + & 5 & & & \\ \text{wt}_4 & & 8 & & 19 & & & 9 & & 4 & & & 4 & & 14 & & & & & & 1 \end{array}$$

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satisfiability **decidable!** (linear Diophantine inequalities)

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1. Check satisfiability for all cycle decompositions of “short” words

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1. Check satisfiability for all cycle decompositions of “short” words
2. “Long words”: one cycle two times  $\implies$  cut

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We show:  $\mathcal{A}_1$  fin-amb  $\implies \mathcal{A}_1 \geq \mathcal{A}_2$  decidable

Cycle Decomposition

$\vec{P}_1 \ x_1 \ \vec{P}_2 \ y_2 \ \vec{P}_2 \ x_3 \ \vec{P}_3 \ y_4 \ \vec{P}_3 \ x_5 \ \vec{P}_2 \ y_2 \ \vec{P}_2 \ x_7 \ \vec{P}_5$

Vectors of weights

$$\begin{array}{rcccccccc} \text{wt}_1 & & 7 & & 13X_1 & & 7 & & 2X_2 & & 12 & & 13X_3 & & 18 \\ \text{wt}_2 & = & 11 & + & 8X_1 & + & 3 & + & 3X_2 & + & 10 & + & 8X_3 & + & 2 \\ \text{wt}_3 & & 4 & + & 6X_1 & + & 1 & + & 15X_2 & + & 9 & + & 6X_3 & + & 5 \\ \text{wt}_4 & & 8 & & 19X_1 & & 9 & & 4X_2 & & 4 & & 19X_3 & & 1 \end{array}$$

$\text{wt}_1, \text{wt}_2, \text{wt}_3 < \text{wt}_4?$  for some choice of  $X_1, X_2, X_3 \in \mathbb{N}$ ?

1. Check satisfiability for all cycle decompositions of “short” words
2. “Long words”: one cycle two times  $\implies$  cut



# THE EQUIVALENCE PROBLEM ON WORDS

We show:  $\mathcal{A}_1$  fin-amb  $\implies \mathcal{A}_1 \geq \mathcal{A}_2$  decidable

Cycle Decomposition

$\vec{P}_1$   $x_1$   $\vec{P}_2$   $y_2$   $\vec{P}_2$   $x_3$   $\vec{P}_3$   $y_4$   $\vec{P}_3$   $x_5$   $\vec{P}_2$   $y_2$   $\vec{P}_2$   $x_7$   $\vec{P}_5$

Vectors of weights

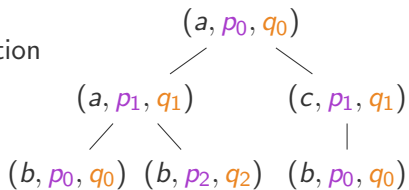
$$\begin{array}{r} \text{wt}_1 \\ \text{wt}_2 \\ \text{wt}_3 \\ \text{wt}_4 \end{array} = \begin{array}{r} 7 \\ 11 \\ 4 \\ 8 \end{array} + \begin{array}{r} 13X_1 \\ 8X_1 \\ 6X_1 \\ 19X_1 \end{array} + \begin{array}{r} 7 \\ 3 \\ 1 \\ 9 \end{array} + \begin{array}{r} 2X_2 \\ 3X_2 \\ 15X_2 \\ 4X_2 \end{array} + \begin{array}{r} 12 \\ 10 \\ 9 \\ 4 \end{array} + \begin{array}{r} 13X_3 \\ 8X_3 \\ 6X_3 \\ 19X_3 \end{array} + \begin{array}{r} 18 \\ 2 \\ 5 \\ 1 \end{array}$$

$\text{wt}_1, \text{wt}_2, \text{wt}_3 < \text{wt}_4?$  for some choice of  $X_1, X_2, X_3 \in \mathbb{N}$ ?

1. Check satisfiability for all cycle decompositions of “short” words
2. “Long words”: one cycle two times  $\implies$  cut

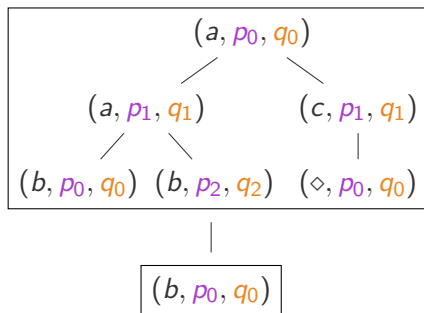
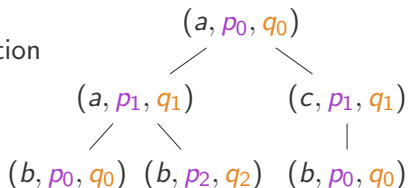
# THE EQUIVALENCE PROBLEM ON TREES

Cycle Decomposition



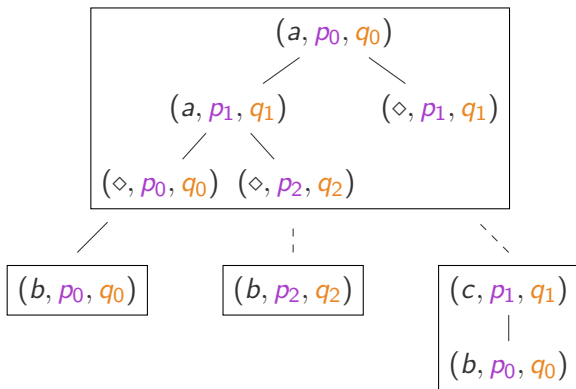
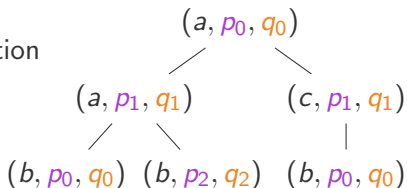
# THE EQUIVALENCE PROBLEM ON TREES

Cycle Decomposition



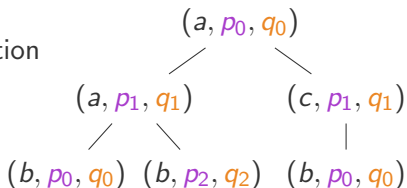
# THE EQUIVALENCE PROBLEM ON TREES

Cycle Decomposition



# THE EQUIVALENCE PROBLEM ON TREES

Cycle Decomposition



Removing Cycles?

