



A DIAMOND for Argumentation

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Motivation

- Dung's abstract argumentation frameworks (AFs) are widely used in argumentation
- They are often criticized for the limitation to simple attacks between arguments
- Abstract Dialectical Frameworks (ADFs) generalize Dung's AFs and allow for arbitrary relations between one argument and a set of *parent* arguments
- DIAMOND (**DI**Alectical **MO**dels **eNc**o**D**ing) is a tool for argumentation researchers to compute extensions of ADFs

Abstract Dialectical Frameworks

Definition. An **abstract dialectical framework** is a tuple $D = (S, L, C)$ where

- S is a set of statements (positions, nodes),
- $L \subseteq S \times S$ is a set of links,
- $C = \{C_s\}_{s \in S}$ is a set of total functions $C_s : 2^{par(s)} \rightarrow \{t, f\}$. C_s is called acceptance condition of s .

Acceptance conditions can be represented as propositional formulae.

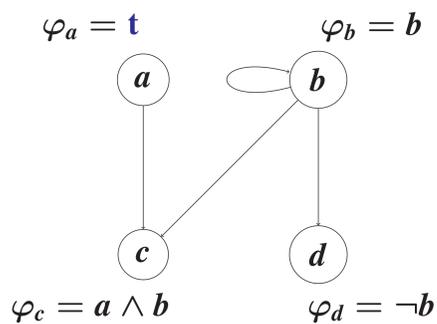


Figure : Example of an ADF

- Semantics are based on three-valued logic
- t , f , and u are partially ordered by \leq_i s.t. $u <_i f$ and $u <_i t$
- **Consensus-Operator**, which assigns $t \sqcap t = t$, $f \sqcap f = f$, and returns u otherwise.

Definition. Let D be an ADF and v be a three-valued interpretation over S . Then the interpretation $\Gamma_D(v)$ is given by

$s \mapsto \sqcap \{w(\varphi_s) \mid w \in [v]_2\}$. Furthermore v is

- **admissible** iff $v \leq_i \Gamma_D(v)$;
- **complete** iff $\Gamma_D(v) = v$, that is, v is a fixpoint of Γ_D ;
- **grounded** iff v is the \leq_i -least fixpoint of Γ_D .

A two-valued interpretation v is a **model of D** iff $\Gamma_D(v) = v$; it is a **stable model of $D = (S, L, C)$** iff v is a model of D and E_v equals the grounded extension of the reduced ADF $D^v = (E_v, L^v, C^v)$, where $L^v = L \cap (E_v \times E_v)$ and for $s \in E_v$ we set $\varphi_s^v = \varphi_s[r/f : v(r) = f]$.

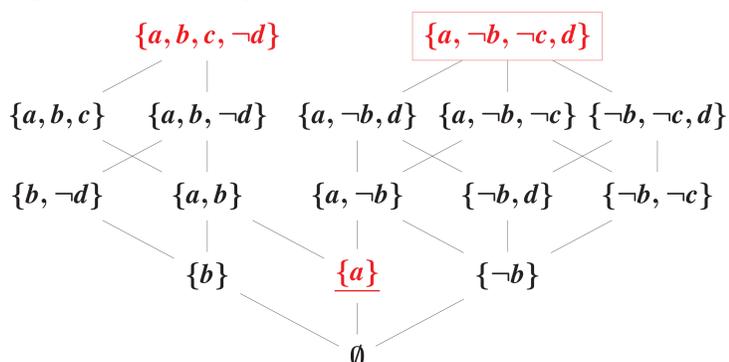


Figure : admissible, complete, grounded, model, and stable model semantics

DIAMOND

- Collection of Answer Set Programming Encodings, utilizing the **Potassco** solvers
- Provides ECLⁱPS^e scripts for input format conversions from **prioritized ADFs** and **propositional formula ADFs** to the **acceptance function** representation
- Command-line interface realized with Python

Instance Representation

An ADF is represented by the predicates

- $s/1$
- $ci/1$
- $co/1$
- $l/2$
- $ci/3$
- $co/3$

where s and l represent the statements and links. ci (resp. co) represents under which accepted parents the acceptance function maps to t (resp. f).

Example

```
s(a). s(b). s(c). s(d).
l(a,c). l(b,b). l(b,c). l(b,d).
ci(a). co(b). ci(b,1,b).
co(c). co(c,1,a). co(c,2,b). ci(c,3,a). ci(c,3,b).
ci(d). co(d,1,b).
```

Implementation Details

- Γ_D will be computed at step(I), based on the three valued interpretation, represented by the predicates $in(X, I)$ and $out(X, I)$
- Γ_D 's result is modeled by $valid(X, I)$, $unsat(X, I)$, and $fp(I)$
- Admissible and complete models can be realized with a **Guess & Check** approach


```
step(0). :- in(S,0), out(S,0).
in(S,0):s(S). :- in(S), not valid(S,0).
out(S,0):s(S). :- out(S), not unsat(S,0).
```
- The computation of the grounded model needs to construct an ordering for iteration until the fixpoint is reached


```
max(I) :- I:=s(S). step(0).
in(S,I+1) :- valid(S,I).
out(S,I+1) :- unsat(S,I).
step(I+1) :- step(I), not max(I).
```
- Stable models combine the **Guess & Check** approach and the grounded model fixpoint-computation

References

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- ▶ Phan Minh Dung. On the Acceptability of Arguments and its Fundamental Role in Nonmonotonic Reasoning, Logic Programming and n-Person Games. *Artificial Intelligence*, 77:321–358, 1995.
- ▶ Martin Gebser, Roland Kaminski, Benjamin Kaufmann, Max Ostrowski, Torsten Schaub, and Marius Thomas Schneider. Potassco: The Potsdam answer set solving collection. *AI Communications*, 24(2):105–124, 2011.