

# Contexts and Preferences in Answer Set Programming

Gerhard Brewka

`brewka@informatik.uni-leipzig.de`

Universität Leipzig

# Motivation 1: preferences

- preferences determine how agents decide and act
- pop up everywhere:

coffee > tea

car > train

relax > work

Chelsea London > Bayern München

Madonna > Britney Spears

marry > don't marry

sleep > listen to talk

# Preferences in AI

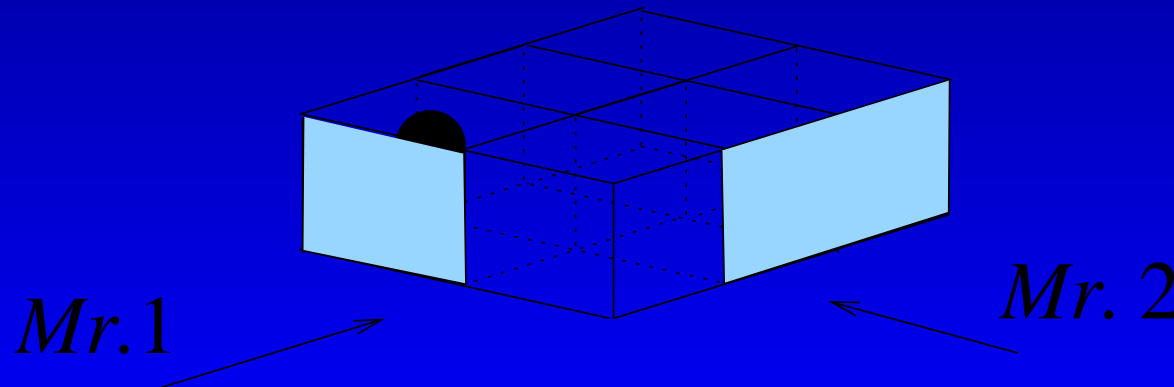
- diagnosis: prefer more plausible hypotheses
- planning/configuration: prefer cheaper plan; satisfy more important constraints
- revision: give up less preferred beliefs
- reasoning about action: prefer fewer unexplained changes
- ontologies: prefer more specific information
- legal/deontic reasoning: prefer more recent law
- linguistics: prefer more important constraints (optimality theory)

# Issues

- how to **represent space of outcomes**  
here: answer sets
- how to **represent preferences**  
here: qualitative orderings
- how to **interpret preferences**  
here: defeasible multi-criteria
- how to **represent dependencies**  
here: rule preconditions specify context

# Motivation 2: contexts

- McCarthy: *we offer no definition of context*
- viewpoints, perspectives, granularity, agents, ...
- (almost) independent units of reasoning
- locality: different languages, reasoning methods
- compatibility: information flow between contexts
- essential for handling complexity, consistency, ...



# Outline

1. Motivation (done)
2. Answer set programming
3. Qualitative preferences
  - Formalisms
  - Applications
4. Contexts
  - Formalism
  - Examples
5. Conclusions

## 2. Answer set programming

# Answer sets

- define semantics for logic programs with strict and default negation (extended LPs)
- rules of the form ( $a, b_i, c_j$  literals):

$$a \leftarrow b_1, \dots, b_n, \text{not } c_1, \dots, \text{not } c_m$$

- AS acceptable set of beliefs based on program
- requirements:
  - closed**: all rules used to generate AS  
if all  $b_i \in \text{AS}$ , no  $c_j \in \text{AS}$ , then  $a \in \text{AS}$
  - grounded**:  $a$  in AS implies derivation of  $a$  from  
rules whose not -preconds are not in AS



# A simple test

To check whether  $S$  is AS of  $P$

- remove  $S$ -defeated rules (not  $L$  in body,  $L \in S$ )
- remove not literals from remaining rules
- check whether  $S =$  closure of reduced program

# A simple test

To check whether  $S$  is AS of  $P$

- remove  $S$ -defeated rules (not  $L$  in body,  $L \in S$ )
- remove not literals from remaining rules
- check whether  $S =$  closure of reduced program

$$a \leftarrow \text{not } b$$
$$b \leftarrow \text{not } c$$

# A simple test

To check whether  $S$  is AS of  $P$

- remove  $S$ -defeated rules (not  $L$  in body,  $L \in S$ )
- remove not literals from remaining rules
- check whether  $S =$  closure of reduced program

$$a \leftarrow \text{not } b$$
$$b \leftarrow \text{not } c$$
$$\{a\}$$

**NO, not closed**

# A simple test

To check whether  $S$  is AS of  $P$

- remove  $S$ -defeated rules (not  $L$  in body,  $L \in S$ )
- remove not literals from remaining rules
- check whether  $S =$  closure of reduced program

$$\begin{array}{l} \cancel{a} \leftarrow \cancel{\text{not } b} \\ b \leftarrow \cancel{\text{not } c} \end{array}$$
$$\{a, b\}$$

**NO, not grounded**

# A simple test

To check whether  $S$  is AS of  $P$

- remove  $S$ -defeated rules (not  $L$  in body,  $L \in S$ )
- remove not literals from remaining rules
- check whether  $S =$  closure of reduced program

$$\begin{array}{l} \cancel{a} \leftarrow \cancel{\text{not } b} \\ b \leftarrow \cancel{\text{not } c} \end{array}$$

$$\{b\}$$

YES, grounded and closed

# Example: graph coloring

Description of graph:

$$\text{node}(v_1), \dots, \text{node}(v_n), \text{edge}(v_i, v_j), \dots$$

Generate: every node needs exactly 1 color

$$\text{col}(X, r) \leftarrow \text{node}(X), \text{not } \text{col}(X, b), \text{not } \text{col}(X, g)$$
$$\text{col}(X, b) \leftarrow \text{node}(X), \text{not } \text{col}(X, r), \text{not } \text{col}(X, g)$$
$$\text{col}(X, g) \leftarrow \text{node}(X), \text{not } \text{col}(X, r), \text{not } \text{col}(X, b)$$

Test: linked nodes cannot have same color

$$\leftarrow \text{edge}(X, Y), \text{col}(X, Z), \text{col}(Y, Z)$$

Each answer set describes a solution!

# Example: graph coloring

Description of graph:

$$\text{node}(v_1), \dots, \text{node}(v_n), \text{edge}(v_i, v_j), \dots$$

Generate: every node needs exactly 1 color

$$\text{col}(X, r) \leftarrow \text{node}(X), \text{not } \text{col}(X, b), \text{not } \text{col}(X, g)$$
$$\text{col}(X, b) \leftarrow \text{node}(X), \text{not } \text{col}(X, r), \text{not } \text{col}(X, g)$$
$$\text{col}(X, g) \leftarrow \text{node}(X), \text{not } \text{col}(X, r), \text{not } \text{col}(X, b)$$

Test: linked nodes cannot have same color

$$f \leftarrow \text{edge}(X, Y), \text{col}(X, Z), \text{col}(Y, Z), \text{not } f$$

Each answer set describes a solution!

# A useful language extension

Bounds on number of satisfied literals:

$$\boxed{L\{a_1, \dots, a_k\}U}$$

Read: at least  $L$  and at most  $U$  of the  $a_i$ s must be true

Allows us to replace 3 color assignment rules with:

$$1\{col(X, r), col(X, b), col(X, g)\}1 \leftarrow node(X)$$

or, if extension of *color* is  $\{r, b, g\}$ :

$$1\{col(X, Y) : color(Y)\}1 \leftarrow node(X)$$



# Why LPs under AS semantics?

- simple yet expressive language
- transitive closure, nonmonotonic rules, ...  
$$flies(X) \leftarrow \text{not } ab(X), bird(X)$$
- simple epistemic distinctions, particularly useful for preference reasoning  
$$safe > \text{not } \neg safe > \neg safe$$
- interesting implementations: dlvs, Smodels, nomore, ASSAT ...
- interesting applications: configuration, diagnosis, planning, reasoning about action, shuttle control, model checking, information integration, ...



# 3. Qualitative Preferences

## 3.1 Formalisms

# Adding preferences to ASP

	rule preference	formula preference
fixed	$(P, <)$ $<$ order on $P$ B-Eiter Delgrande-Schaub ...	$(P, <)$ $<$ order on $Lit$ Sakama-Inoue Foo-Zhang ...
conditional	$<$ predicate in $P$ applied to rules B-Eiter Delgrande-Schaub ...	ordered disjunction ASO programs B-Niemelä-Syrjänen B-N-Truszczyński ...

# Ordered disjunction

LPOD: finite set of rules of the form:

$$C_1 \times \dots \times C_n \leftarrow \text{body}$$

*if body then some  $C_j$  must be true, preferably  $C_1$ , if impossible then  $C_2$ , if impossible  $C_3$ , etc.*

- answer sets defined through split programs
- satisfy rules to different degrees, depending on best satisfied head literal
- use degrees to define global preference relation on answer sets
- different options how to do this

# Preferences among answer sets

How to generate global preference ordering from satisfaction degrees?

Many options, for instance:

$P^i(S)$  =  $P$ -rules  $i$ -satisfied in  $S$ .  $S_1 > S_2$  iff

- some rule has better satisfaction degree in  $S_1$  and no rule better degree in  $S_2$ ,
- at smallest degree  $i$  with  $P^i(S_1) \neq P^i(S_2)$ ,  $S_1$  satisfies superset of rules satisfied in  $S_2$ ,
- at smallest degree  $i$  with  $|P^i(S_1)| \neq |P^i(S_2)|$ ,  $S_1$  satisfies more rules than  $S_2$ .

# Prioritized graph coloring

$$\begin{aligned} &col(X, r) \times col(X, b) \times col(X, g) \leftarrow node(X) \\ &\leftarrow col(X, C), col(Y, C), edge(X, Y) \end{aligned}$$

$M$  preferred over  $M'$  if

- par* at least 1 node has nicer color in  $M$  than in  $M'$ ,  
no node less preferred color.
- incl* nodes red in  $M$  superset of nodes red in  $M'$ , or  
same nodes red in  $M$  and  $M'$  and nodes blue in  
 $M$  superset of nodes blue in  $M'$ .
- card* more nodes red in  $M$  than in  $M'$ , or as many  
nodes red in  $M$  as in  $M'$  and more blue in  $M$ .

# The ASO approach

- decoupled approach to answer set optimization
- logic program  $G$  generates answer sets
- preference program  $P$  used to compare them
- preference program set of rules

$$C_1 > \dots > C_k \leftarrow \text{body}$$

- $C_i$  boolean combination built using  $\vee$ ,  $\wedge$ ,  $\neg$ , not
- rule satisfaction and combination as for LPODs

# LPODs vs. ASO

- ASO: arbitrary generating programs, no implicit generation of options, general preferences:  
combinations of properties preferred over others:

$$a > (b \wedge c) > d \leftarrow f$$

equally preferred options:

$$a > (b \vee c) > \text{not } d \leftarrow g$$

- LPODs: compact and readable representations



## 3.2 Applications

# Configuration

- often represented as AND/OR trees
- simple representation with Smodels cardinalities:

$4\{starter, main, dessert, drink\}4 \leftarrow dinner$

$1\{soup, salad\}1 \leftarrow starter$

$1\{fish, beef, lasagne\}1 \leftarrow main$

- add case description and preferences, e.g.

$fish \vee beef > lasagne$

$beer > wine \leftarrow beef$

$wine > beer \leftarrow \text{not } beef$

- preferred answer sets: optimal configurations

# Abductive diagnosis

*H* : measles, flu, migraine

*O* : headache, fever

*K* : fever  $\leftarrow$  measles

red-spots  $\leftarrow$  measles

headache  $\leftarrow$  migraine

nausea  $\leftarrow$  migraine

fever  $\leftarrow$  flu

headache  $\leftarrow$  flu

# Abductive diagnosis

$H$  : *measles, flu, migraine*

$O$  : *headache, fever*

$K$  : *fever*  $\leftarrow$  *measles*

*red-spots*  $\leftarrow$  *measles*

*headache*  $\leftarrow$  *migraine*

*nausea*  $\leftarrow$  *migraine*

*fever*  $\leftarrow$  *flu*

*headache*  $\leftarrow$  *flu*

$\neg$ *measles*  $\times$  *measles*

$\neg$ *flu*  $\times$  *flu*

$\neg$ *migraine*  $\times$  *migraine*

# Abductive diagnosis

$H$  : *measles, flu, migraine*

$O$  : *headache, fever*

$K$  : *fever*  $\leftarrow$  *measles*

*red-spots*  $\leftarrow$  *measles*

*headache*  $\leftarrow$  *migraine*

*nausea*  $\leftarrow$  *migraine*

*fever*  $\leftarrow$  *flu*

*headache*  $\leftarrow$  *flu*

$\neg$ *measles*  $\times$  *measles*

$\neg$ *flu*  $\times$  *flu*

$\neg$ *migraine*  $\times$  *migraine*

$\leftarrow$  not *headache*

$\leftarrow$  not *fever*

# Abductive diagnosis

$H$  : *measles, flu, migraine*

$O$  : *headache, fever*

$K$  : *fever*  $\leftarrow$  *measles*

*red-spots*  $\leftarrow$  *measles*

*headache*  $\leftarrow$  *migraine*

*nausea*  $\leftarrow$  *migraine*

*fever*  $\leftarrow$  *flu*

*headache*  $\leftarrow$  *flu*

$\neg$ *measles*  $\times$  *measles*

$\neg$ *flu*  $\times$  *flu*

$\neg$ *migraine*  $\times$  *migraine*

$\leftarrow$  not *headache*

$\leftarrow$  not *fever*

inclusion preferred:  $\{migraine, measles\}, \{flu\}$

# Inconsistency handling

- program  $P$ , possibly inconsistent; consistency restoring rules  $R$
- names  $N_P$  and  $N_R$  for rules in  $P$  and  $R$
- generate weakening of  $P \cup R$  by replacing

$$\boxed{head \leftarrow body} \quad \text{with} \quad \boxed{head \leftarrow body, r_i}$$

where  $r_i$  rule's name

- add  $\{r \times \neg r \mid r \in N_P\} \cup \{\neg r \times r \mid r \in N_R\}$
- minimal set of  $P$ -rules switched off, minimal set of  $R$ -rules switched on

# Game theory

## Prisoners' dilemma

	Coop.	Defect
Coop.	3,3	0,5
Defect	5,0	1,1



# Game theory

## Prisoners' dilemma

	Coop.	Defect
Coop.	3,3	0,5
Defect	5,0	1,1

Player 1:

$$D_1 \times C_1 \leftarrow C_2$$

$$D_1 \times C_1 \leftarrow D_2$$

Player 2:

$$D_2 \times C_2 \leftarrow C_1$$

$$D_2 \times C_2 \leftarrow D_1$$

Move clause:  $1\{C_1, D_1\}1$

# Game theory

## Prisoners' dilemma

	Coop.	Defect
Coop.	3,3	0,5
Defect	5,0	1,1

Player 1:

$$D_1 \times C_1 \leftarrow C_2$$

$$D_1 \times C_1 \leftarrow D_2$$

Player 2:

$$D_2 \times C_2 \leftarrow C_1$$

$$D_2 \times C_2 \leftarrow D_1$$

Move clause:  $1\{C_1, D_1\}1$

Preferred answer set = **Nash equilibrium**

# Related issues

- **meta-preferences:** one preference rule/ordered disjunction more important than another
- **preference description language:** combines different preference strategies; integrates qualitative with quantitative methods
- **implementation:** *generate and improve* method; iterative calls to answer set solver generate sequence of strictly improving answer sets
- **integration with CP-nets:** combines graph based methods with flexibility of ASO preferences



# 4. Contexts

## 4.1 Formalism

# Contexts: the Trento school

Multi-context system

$$(\{T_i\}, Delta_{br})$$

each  $T_i = (L_i, \Omega_i, \Delta_i)$  formal system,  $Delta_{br}$  bridge rules using labeled formulas  $c : p$  with  $p \in L_c$ .

- semantics: local models + compatibility
- information flow across contexts via bridge rules
- reasoning within/across contexts monotonic
- exception (Roelofsen/Serafini 05) has problems

# Contextual default logic

Contextual default theory  $((D_1, W_1), \dots, (D_n, W_n))$ .

$(D_i, W_i)$  default theory, default rules in  $D_i$  possibly refer to other contexts.

$\Gamma(S_1, \dots, S_n)$  minimal tuple  $(S'_1, \dots, S'_n)$  with:

1.  $W_i \subseteq S'_i$ ,
2.  $S'_i$  deductively closed (over  $L_i$ ), and
3.  $(c_1 : p_1), \dots, (c_t : p_t) : (c_{t+1} : q_1), \dots, (c_{t+k} : q_k) / r \in D_i$  and for all  $i, j$ :  $p_i \in S'_{c_i}$  and  $\neg q_j \notin S'_{c_{t+j}}$ , then  $r \in S'_i$ .

Extension fixpoint of  $\Gamma$ .

# Contextual logic programming

- Contextual LP system:  $C = (P_1, \dots, P_n)$ .
- $P_i$  contextual LP: labeled literals  $(c : l)$  allowed in rule bodies ( $c$  context,  $l$  literal in  $c$ 's language).
- Minimal context model of definite  $C$  (no **not**): smallest  $(S_1, \dots, S_n)$  with:
  1.  $a \in S_i$  if  $a \leftarrow (c_1 : b_1), \dots, (c_k : b_k) \in P_i$ ,  
 $b_1 \in S_{c_1}, \dots, b_k \in S_{c_k}$ ,
  2.  $S_i = Lit_i$  if  $S_i$  has complementary literals.

# Answer sets

- $C = (P_1, \dots, P_n)$  arbitrary system,  
 $S = (S_1, \dots, S_n)$ ,  $S_i$  literals in  $P_i$ 's language.
- $C^S$  obtained from  $C$  by
  1. deleting rules with literal **not**  $(c:l)$  s.t.  $l \in S_c$ ,
  2. deleting **not** literals from remaining rules.
- $C^S$  definite, has minimal context model  $M_{C^S}$ .
- $S$  answer set iff  $S = M_{C^S}$ .

Also possible to define skeptical inference à la WFS



## 4.2 Examples

# Abstraction

Car domain,  $P_1$  abstraction of  $P_2$ ,  $P_1$  may contain

$expensive(X) \leftarrow 2 : price(X, Y), Y \geq 30000$

$sportive(X) \leftarrow 2 : speed(X, Y), 180 < Y,$   
**not**  $bad-accel(X)$

$bad-accel(X) \leftarrow 2 : 0-to-100(X, Y), Y > 10$

Assume now intelligent grounder instantiating context, proposition and domain variables correctly.

# Information fusion

believe  $p$  if someone does and nobody believes  $\neg p$ :

$$\begin{aligned} P &\leftarrow (C : P), \mathbf{not} \text{ rej}(P) \\ \text{rej}(P) &\leftarrow (C : \neg P) \end{aligned}$$

believe  $p$  if someone you trust does and nobody you trust believes  $\neg p$ :

$$\begin{aligned} P &\leftarrow (C : P), \text{ trusted}(C), \mathbf{not} \text{ rej}(P) \\ \text{rej}(P) &\leftarrow (C : \neg P), \text{ trusted}(C) \end{aligned}$$

believe  $p$  if majority does:

$$P \leftarrow N\{(C : P) : \text{con}(C)\}N, N > n/2$$

# Information fusion, ctd.

total preference order via numbering:  $1 < 2 < 3 \dots$

$$\begin{aligned} P &\leftarrow (C : P), \mathbf{not} \text{rej}(C, P) \\ \text{rej}(C, P) &\leftarrow (C : P), (C' : \neg P), C < C' \end{aligned}$$

partial preference order via predicate  $\prec$ , sceptical

$$\begin{aligned} P &\leftarrow \text{acc}(P), \mathbf{not} \text{acc}(\neg P) \\ \text{acc}(P) &\leftarrow (C : P), \mathbf{not} \text{rej}(C, P) \\ \text{rej}(C, P) &\leftarrow (C : P), (C' : \neg P), C \prec C' \end{aligned}$$

# Prisoners' dilemma II

2-context system  $(P_1, P_2)$  with  $P_1$ :

$choose(d) \leftarrow \mathbf{not} choose(c)$

$choose(c) \leftarrow \mathbf{not} choose(d)$

$best(d) \leftarrow (2 : choose(c))$

$best(d) \leftarrow (2 : choose(d))$

$\leftarrow choose(X), \mathbf{not} best(X)$

$P_2 = P_1$  with 2 replaced by 1.

Answer set = Nash equilibrium:

$(\{choose(d), best(d)\}, \{choose(d), best(d)\})$

# Voting

Simple majority vote:

$$votes(X, N) \leftarrow N\{(C : best(X)) : con(C)\}N, cnd(X)$$

$$wins(X) \leftarrow \mathbf{not} \neg wins(X)$$

$$\neg wins(X) \leftarrow votes(X, N), votes(Y, M), M > N$$

Condorcet rule:

$$beats(X, Y, N) \leftarrow N\{(C : beats(X, Y)) : con(C)\}N, \\ cnd(X), cnd(Y)$$

$$wins(X) \leftarrow \mathbf{not} \neg wins(X)$$

$$\neg wins(X) \leftarrow beats(X, Y, N), beats(Y, X, M), \\ M \geq N$$

# 5. Conclusions

# What has been achieved?

- ASP a promising declarative paradigm
- simple yet expressive, interesting applications, interesting solvers
- adding preferences has great potential
- presented two approaches and possible applications
- showed how to add simple notion of context
- missing: current context, switching context, identification of most adequate context, ...



# Future work

Better integration of contexts and preferences

- use prioritized logic programs within context
- use priorities among bridge rules
- modify priorities in a context using bridge rules

Thanks to:

I. Niemelä, M. Truszczyński, F. Roelofsen, L. Serafini;  
you for listening

# References

1. G. Brewka, Logic programs with ordered disjunction, Proc. AAI-02, Edmonton, 2002
2. G. Brewka, I. Niemelä, M. Truszczynski: Answer set optimization, Proc. IJCAI-03, Acapulco, 2003
3. G. Brewka, S. Benferhat, D. Le Berre: Qualitative choice logic, Artificial Intelligence 157(1-2), Special Issue on Nonmonotonic Reasoning, 2004, 203-237
4. G. Brewka, I. Niemelä, T. Syrjänen: Logic programs with ordered disjunction, Computational Intelligence 20 (2), 2004, 335-357
5. G. Brewka: A rank based description language for qualitative preferences, Proc. ECAI-04, Valencia, 2004, 303-307 (also in: Proc. NMR 2004, 88-93)
6. G. Brewka: Complex preferences for answer set optimization, Proc. KR-04, Whistler, Canada, AAAI Press, 2004, 213-223
7. G. Brewka, I. Niemelä, M. Truszczynski: Prioritized component systems, Proc. AAI-05, Pittsburgh 2005
8. G. Brewka, J. Dix: Knowledge Representation with Logic Programs, in: D. Gabbay and F. Guentner (eds.), Handbook of Philosophical Logic, Vol. 12, Springer, 2005, 1-85
9. G. Brewka: Answer Sets and Qualitative Decision Making, Synthese 146 (2005), 171-181
10. G. Brewka, F. Roelofsen, L. Serafini: Contextual Default Reasoning, IJCAI-07