

Constraint Satisfaction Problems

**A Quick Overview
(based on AIMA book slides)**

Constraint satisfaction problems

- **What is a CSP?**
 - *Finite set of variables* V_1, V_2, \dots, V_n
 - *Nonempty domain of possible values* for each variable $D_{V_1}, D_{V_2}, \dots, D_{V_n}$
 - *Finite set of constraints* C_1, C_2, \dots, C_m
 - Each *constraint* C_i limits the values that variables can take, e.g., $V_1 \neq V_2$
- A **state** is an **assignment** of values to some or all variables.
- **Consistent assignment:** assignment does not violate the constraints.

Constraint satisfaction problems

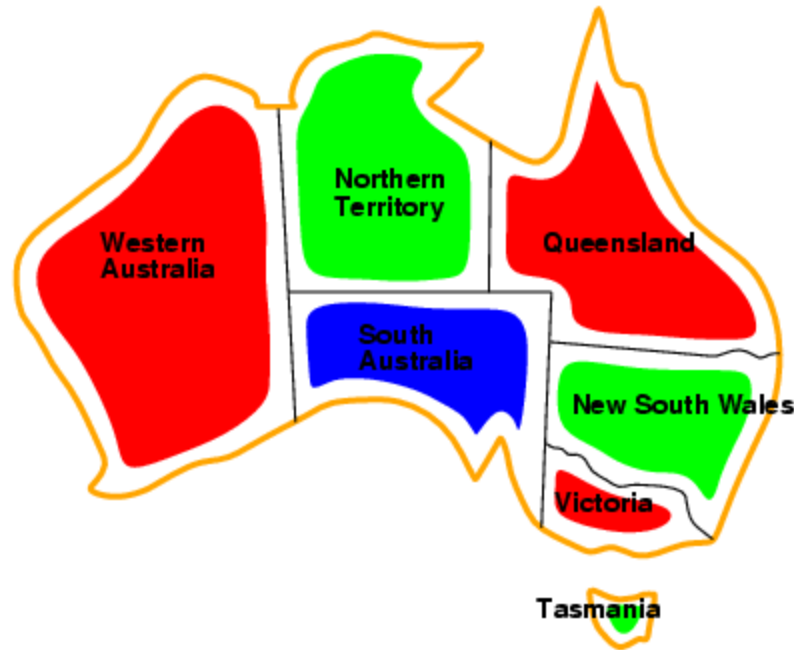
- An assignment is *complete* when every variable has a value.
- A *solution* to a CSP is a complete assignment that satisfies all constraints.
- Some CSPs require a solution that maximizes an *objective function*.
- **Applications:**
 - Scheduling the Hubble Space Telescope,
 - Floor planning for VLSI,
 - Map coloring,
 - Cryptography

Example: Map-Coloring



- **Variables:** *WA, NT, Q, NSW, V, SA, T*
- **Domains:** $D_i = \{\text{red, green, blue}\}$
- **Constraints:** adjacent regions must have different colors
 - e.g., $WA \neq NT$
 - So (WA, NT) must be in $\{(\text{red, green}), (\text{red, blue}), (\text{green, red}), \dots\}$

Example: Map-Coloring

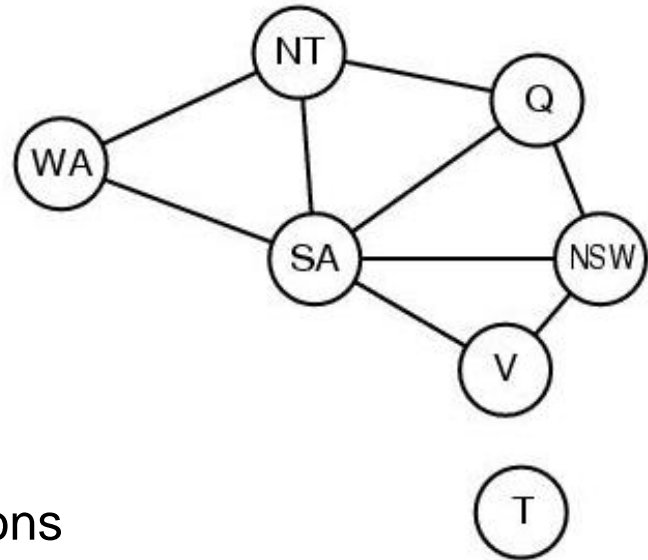


Solutions are **complete** and **consistent** assignments,

- e.g., WA = red, NT = green, Q = red, NSW = green, V = red, SA = blue, T = green

Constraint graph

- **Binary CSP:** each constraint relates two variables
- **Constraint graph:**
 - *nodes* are variables
 - *arcs* are constraints
- **CSP benefits**
 - Standard representation pattern
 - Generic goal and successor functions
 - Generic heuristics (no domain specific expertise).
- **Graph can be used to simplify search.**
 - e.g. Tasmania is an independent subproblem.



Varieties of CSPs

- **Discrete variables**

- finite domains:
 - n variables, domain size $d \rightarrow O(d^n)$ complete assignments
 - e.g., Boolean CSPs, includes Boolean satisfiability (NP-complete)
- infinite domains:
 - integers, strings, etc.
 - e.g., job scheduling, variables are start/end days for each job
 - need a constraint language, e.g., $StartJob_1 + 5 \leq StartJob_3$

- **Continuous variables**

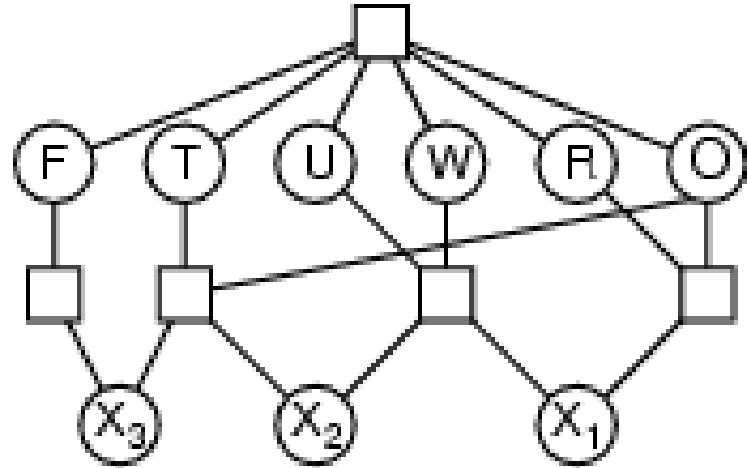
- e.g., start/end times for Hubble Space Telescope observations
- linear constraints solvable in polynomial time by linear programming

Varieties of constraints

- ***Unary*** constraints involve a single variable,
 - e.g., SA \neq green
- ***Binary*** constraints involve pairs of variables,
 - e.g., SA \neq WA
- ***Higher-order*** constraints involve 3 or more variables
 - e.g., cryptarithmic column constraints
- ***Preference*** (soft constraints) e.g. *red is better than green* can be represented by a cost for each variable assignment => **Constrained optimization** problems.

Example: Cryptarithmic

$$\begin{array}{r}
 \text{ T W O} \\
 + \text{ T W O} \\
 \hline
 \text{ F O U R}
 \end{array}$$



- **Variables:**
 $FTUWR O X_1 X_2 X_3$
- **Domain:** $\{0,1,2,3,4,5,6,7,8,9\}$
- **Constraints:** *Alldiff* (F,T,U,W,R,O)
 - $O + O = R + 10 \cdot X_1$
 - $X_1 + W + W = U + 10 \cdot X_2$
 - $X_2 + T + T = O + 10 \cdot X_3$
 - $X_3 = F, T \neq 0, F \neq 0$

CSP as a standard search problem

- **A CSP can easily be expressed as a standard search problem.**
 - *Initial State*: the empty assignment {}.
 - *Operators*: Assign value to unassigned variable provided that there is no conflict.
 - *Goal test*: assignment consistent and complete.
 - *Path cost*: constant cost for every step.
 - Solution is found at depth n , for n variables
 - Hence depth first search can be used

Backtracking search

- Variable assignments are *commutative*,
 - Eg [*WA = red then NT = green*]
equivalent to [*NT = green then WA = red*]
- Only need to consider assignments to a single variable at each node
 - $b = d$ and there are d^n leaves
- Depth-first search for CSPs with single-variable assignments is called *backtracking* search
- Backtracking search basic uninformed algorithm for CSPs
- Can solve *n*-queens for $n \approx 25$

Backtracking search

```
function BACKTRACKING-SEARCH(csp) % returns a solution or failure  
  return RECURSIVE-BACKTRACKING({} , csp)
```

```
function RECURSIVE-BACKTRACKING(assignment, csp) % returns a solution or failure  
  if assignment is complete then return assignment  
  var ← SELECT-UNASSIGNED-VARIABLE(VARIABLES[csp],assignment,csp)  
  for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do  
    if value is consistent with assignment according to CONSTRAINTS[csp]  
    then  
      add {var=value} to assignment  
      result ← RECURSIVE-BACKTRACKING(assignment, csp)  
      if result ≠ failure then return result  
      remove {var=value} from assignment  
  return failure
```

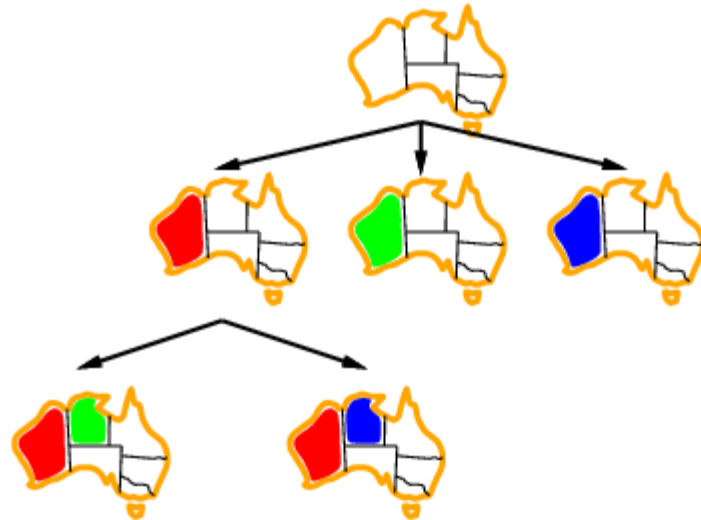
Backtracking example



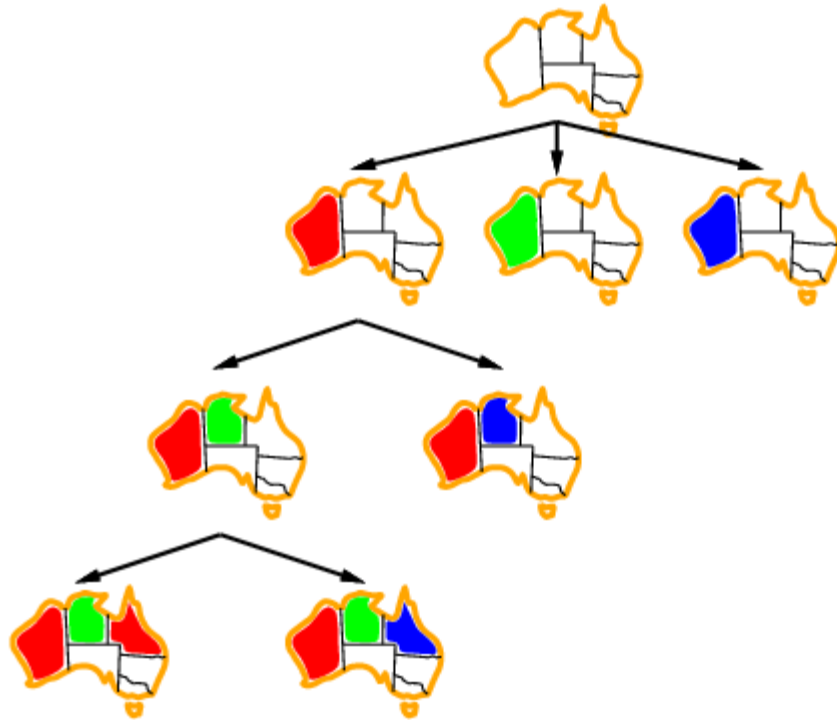
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Backtracking example



Backtracking example



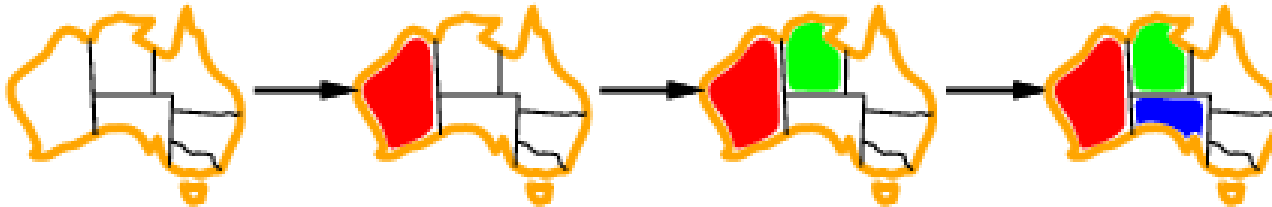
Improving backtracking efficiency

***General-purpose* methods can give huge speed gains:**

- Which variable should be assigned next?
- In what order should its values be tried?
- Can we detect inevitable failure early?

Most constrained variable

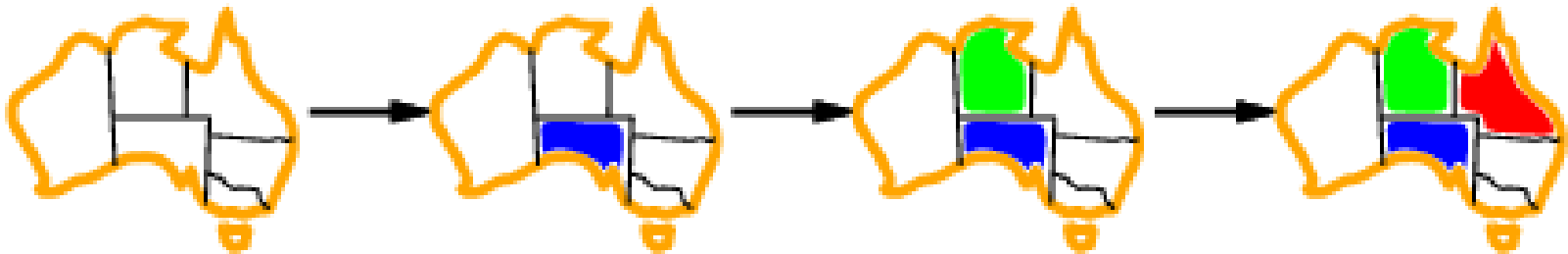
- **Most constrained variable:**
choose the variable with the fewest legal values



- a.k.a. *minimum remaining values (MRV)* heuristic

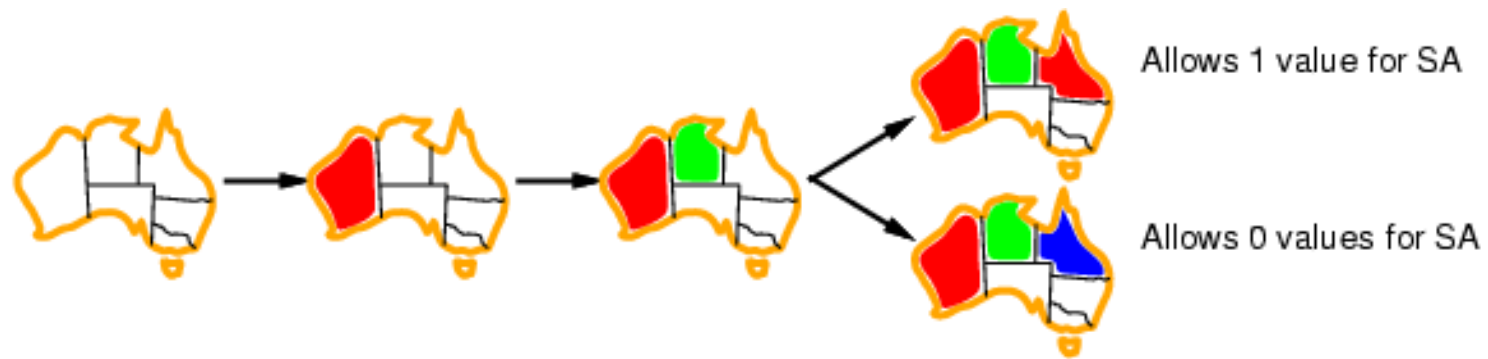
Most constraining variable

- **Tie-breaker among most constrained variables**
- **Most constraining variable:**
 - choose the variable with the most constraints on remaining variables



Least constraining value

- **Given a variable, choose the least constraining value:**
 - the one that rules out the fewest values in the remaining variables



- **Combining these heuristics makes 1000 queens feasible**

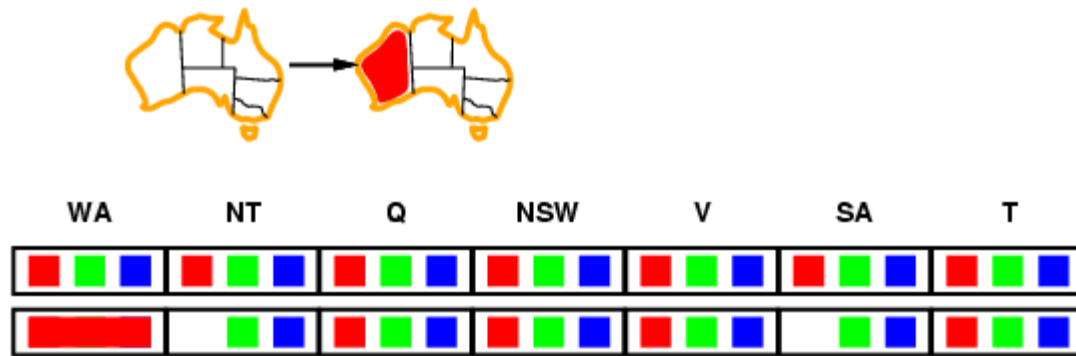
Forward Checking

- **Idea:**
 - Keep track of remaining legal values for unassigned variables
 - Terminate search when any variable has no legal values



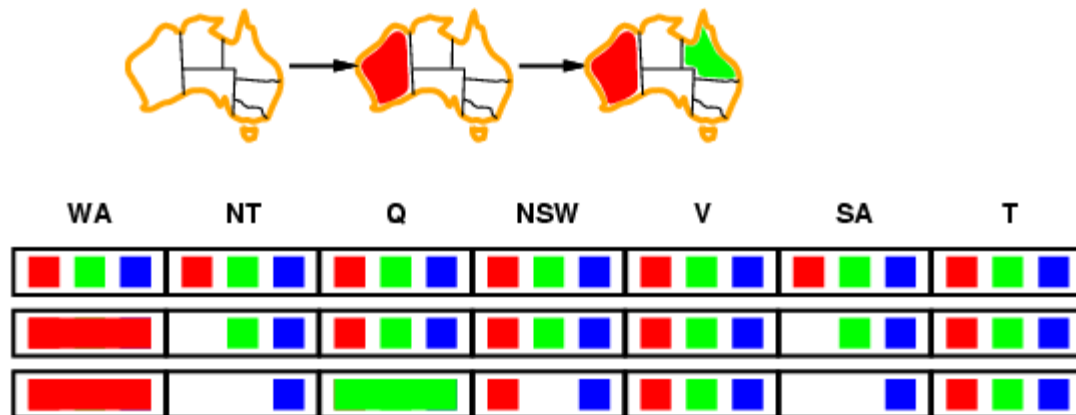
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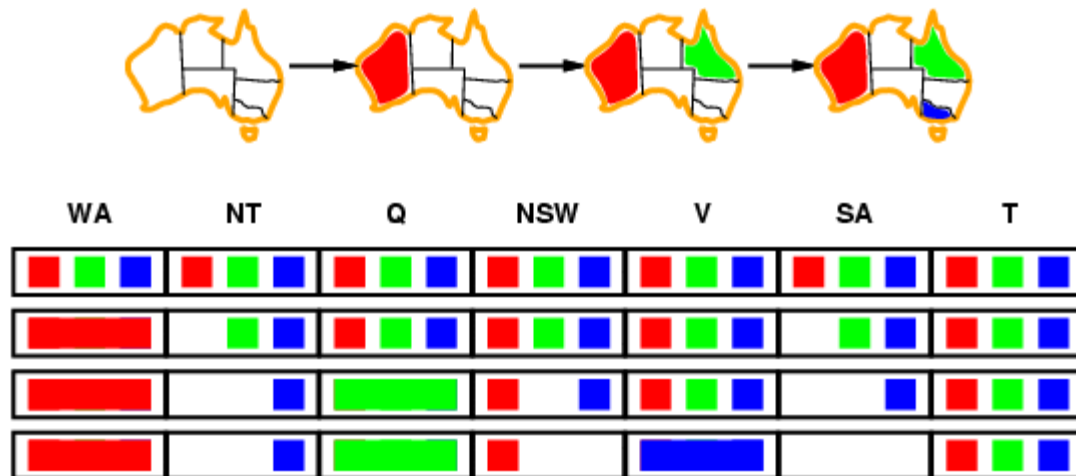
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Forward checking

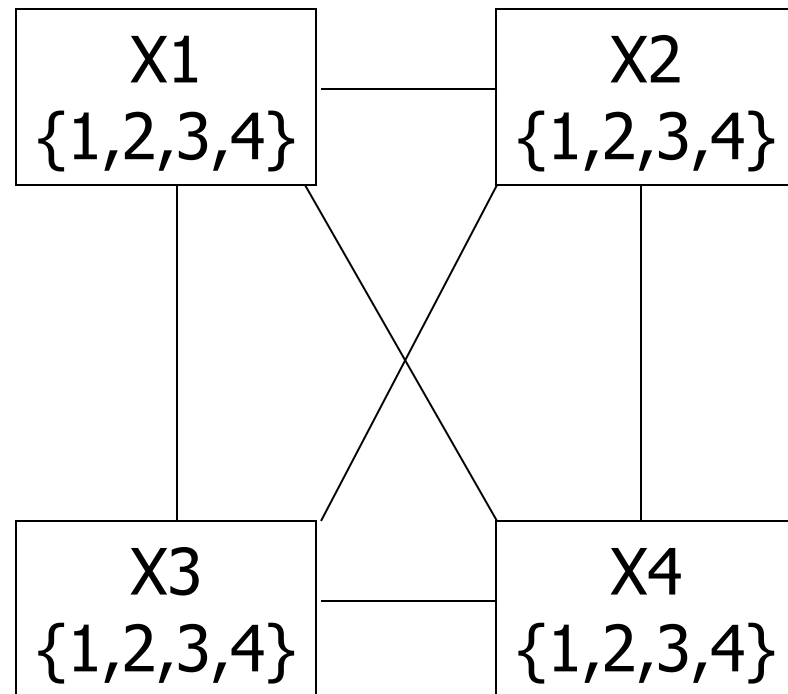
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- **No more value for SA: backtrack**

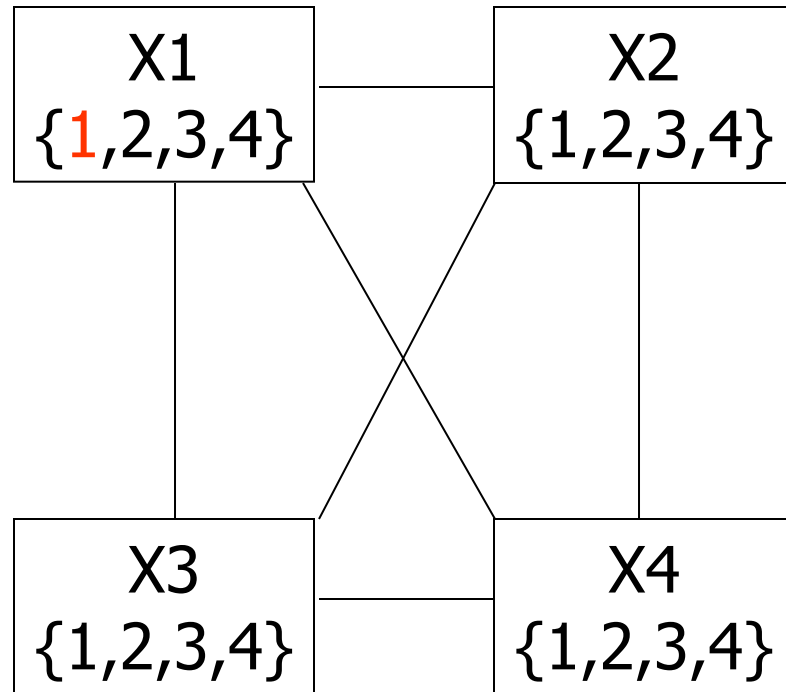
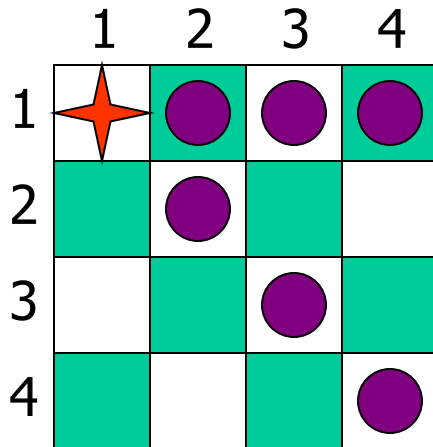
Example: 4-Queens Problem

	1	2	3	4
1				
2				
3				
4				

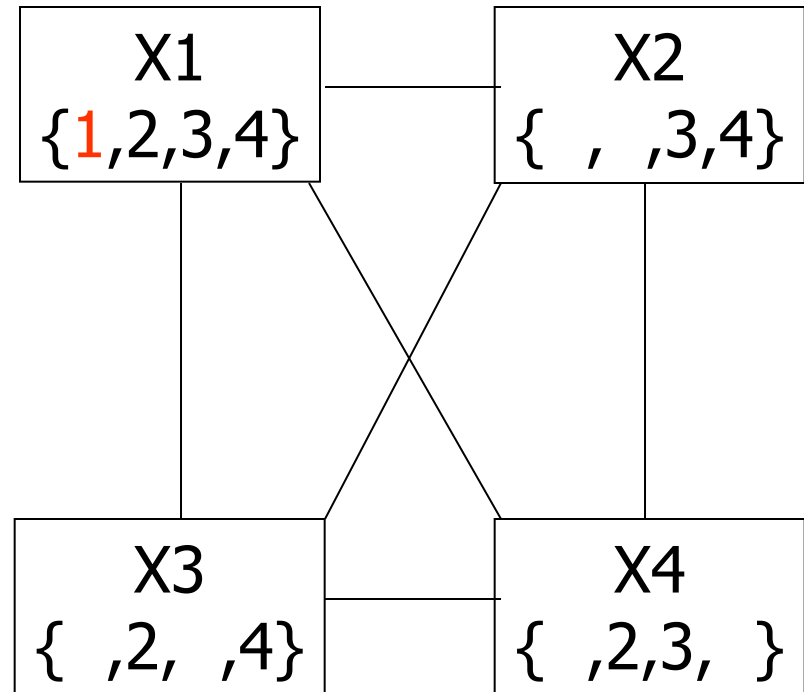
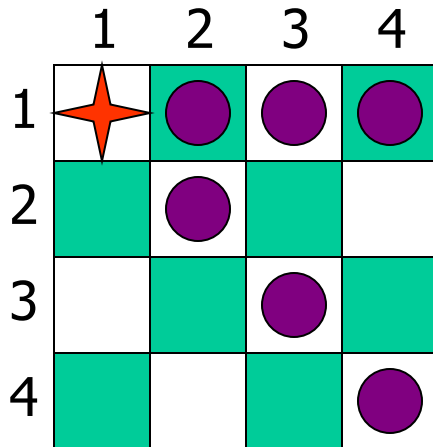


[4-Queens slides copied from B.J. Dorr]

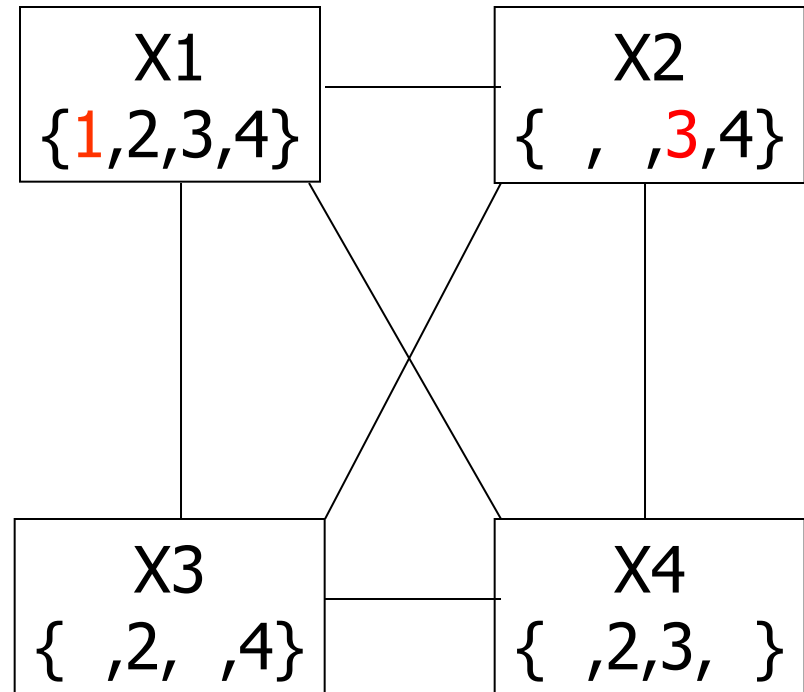
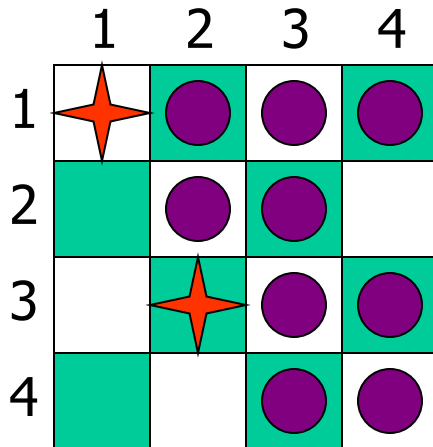
Example: 4-Queens Problem



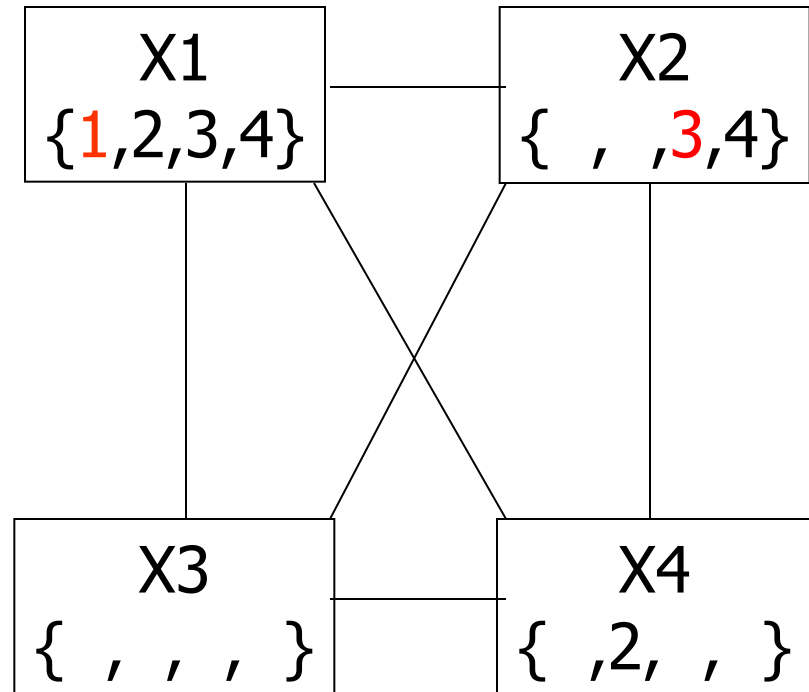
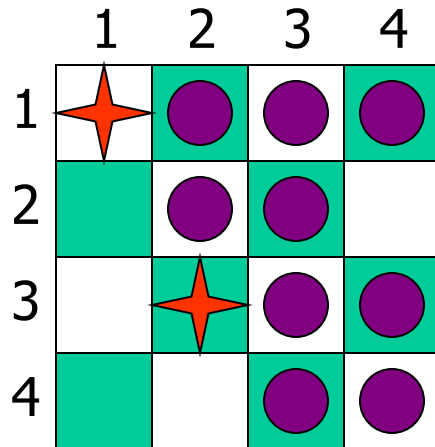
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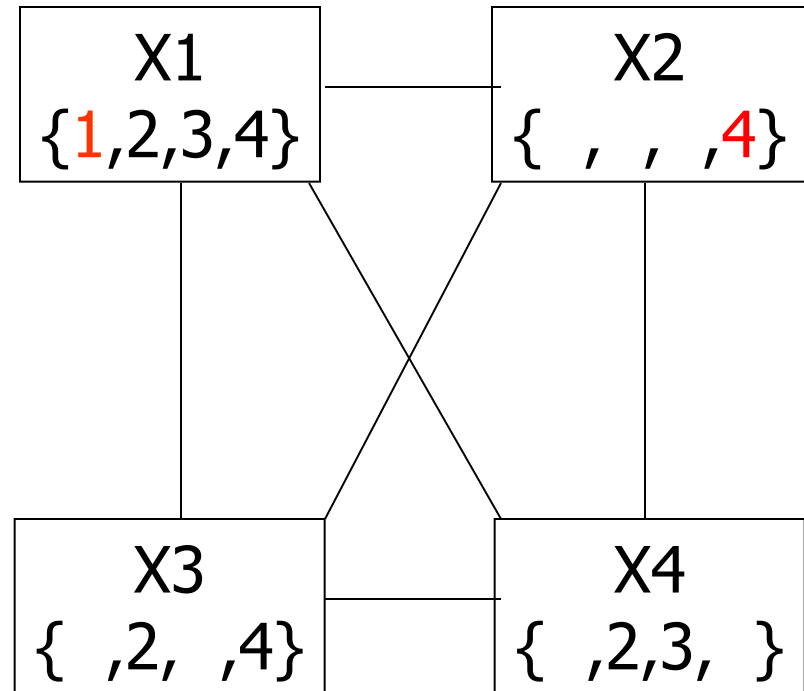
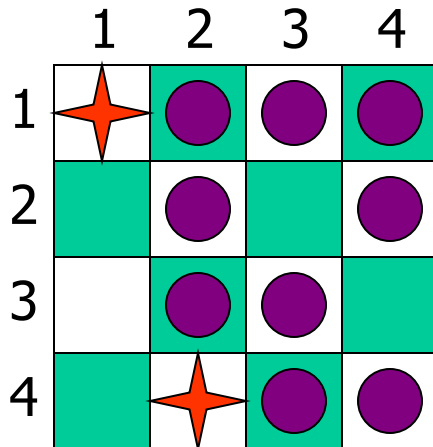
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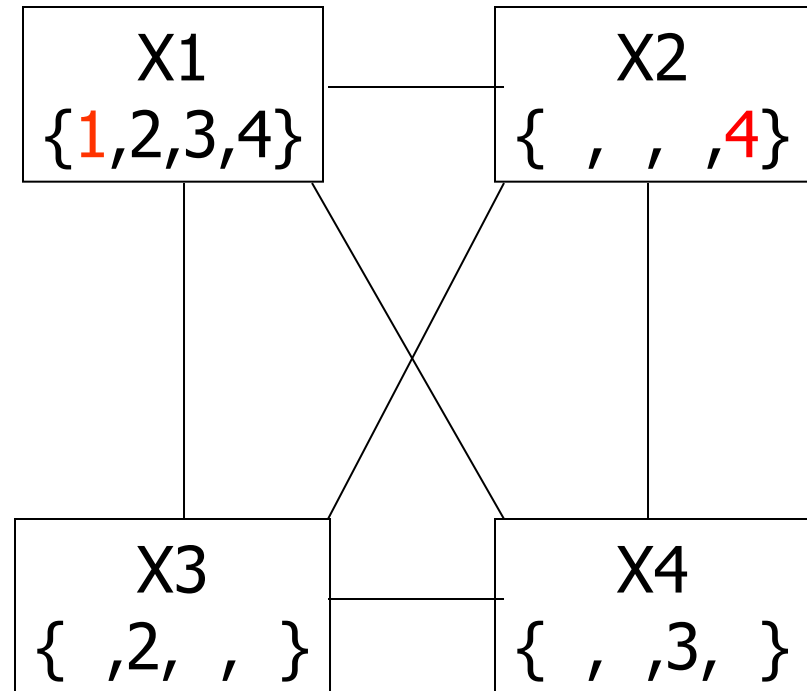
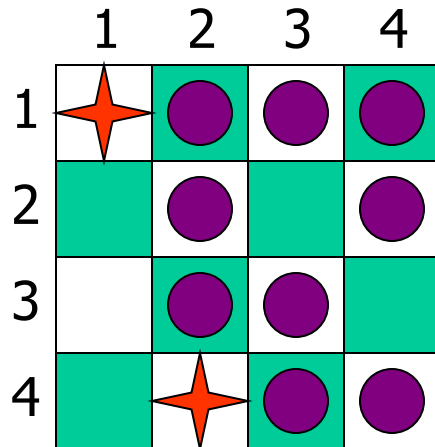
Example: 4-Queens Problem



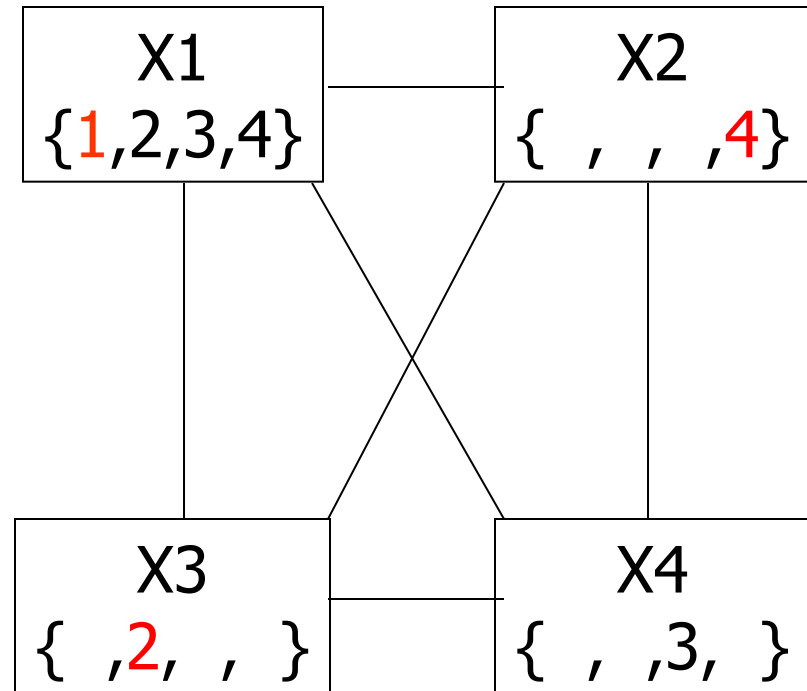
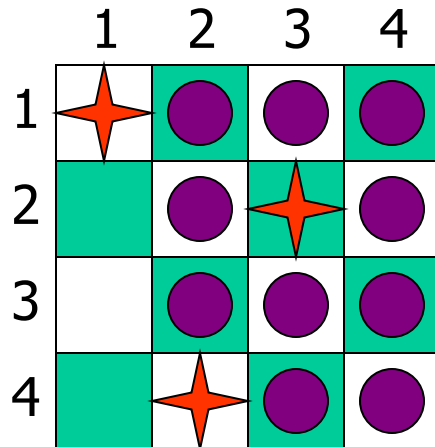
Example: 4-Queens Problem



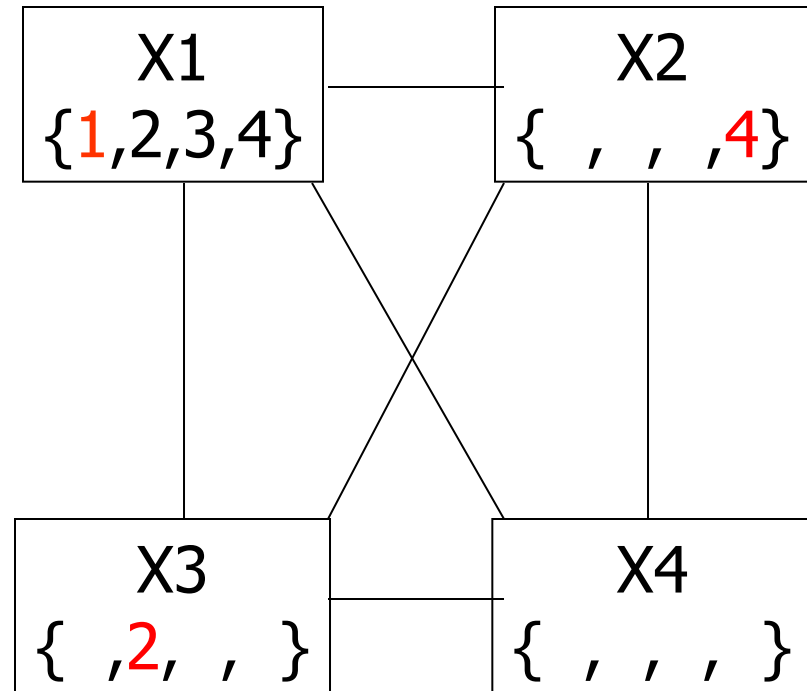
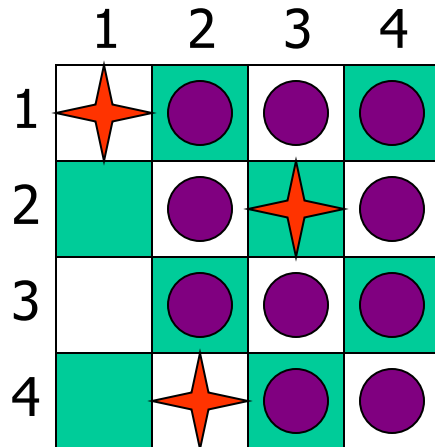
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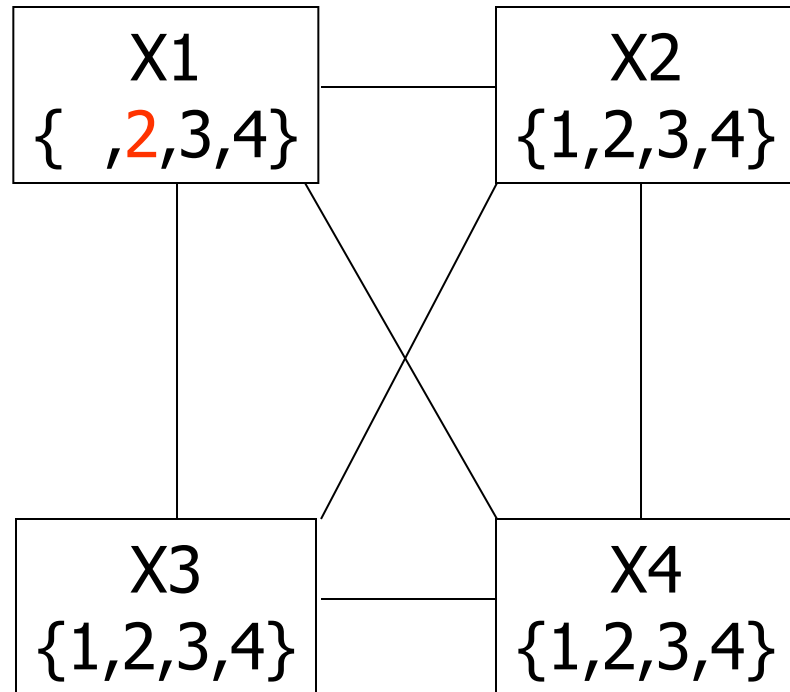
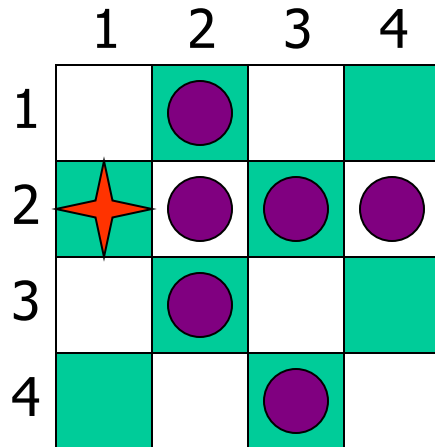
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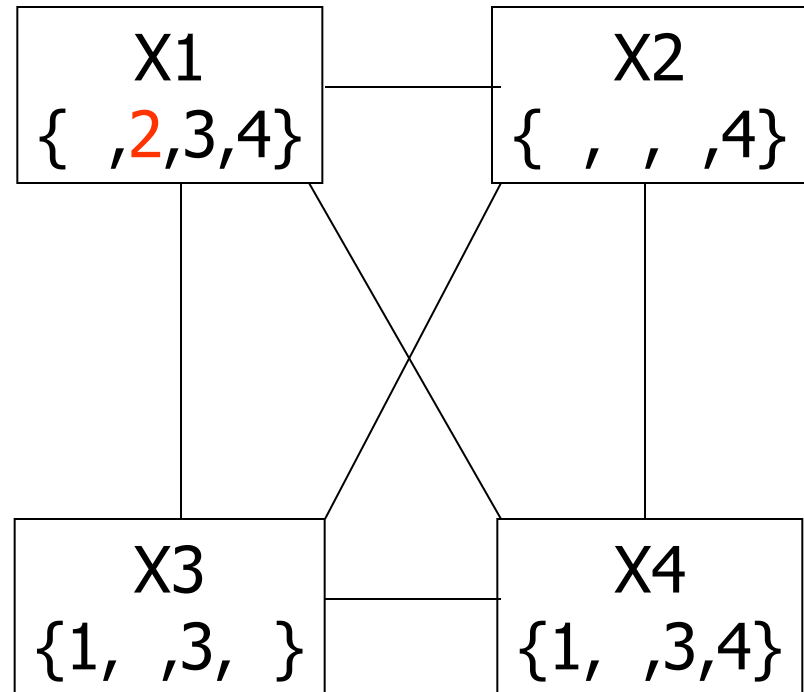
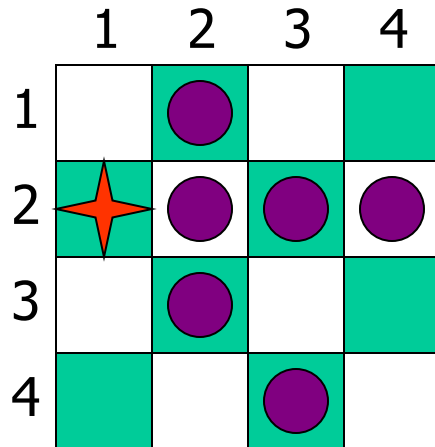
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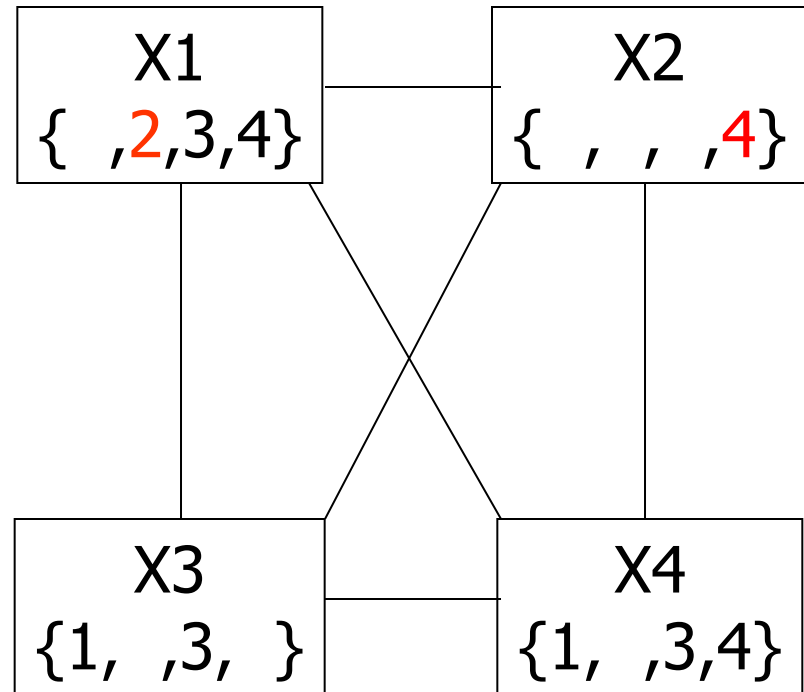
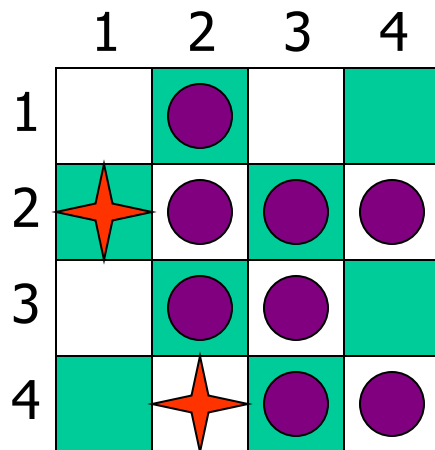
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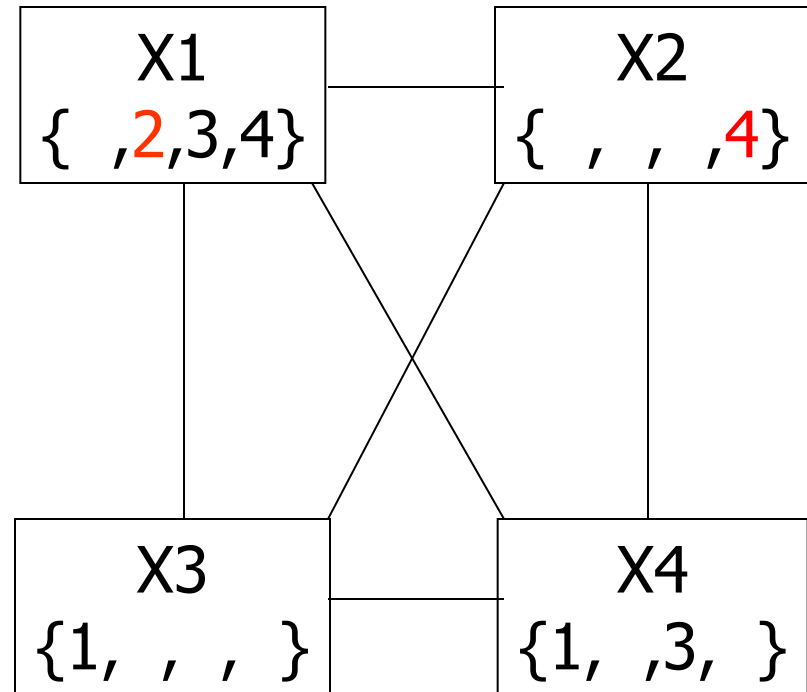
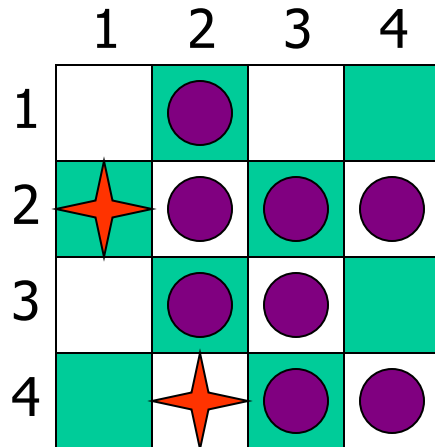
Example: 4-Queens Problem



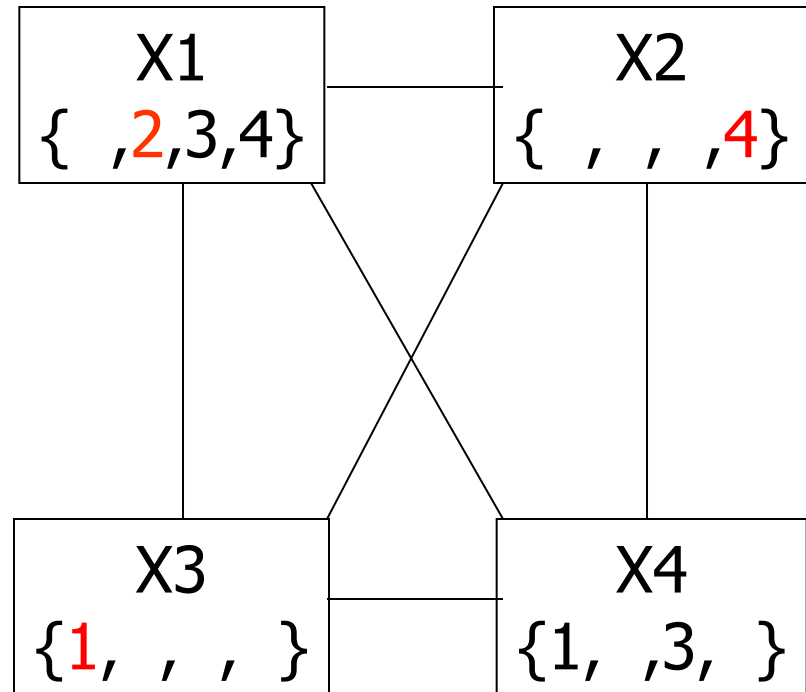
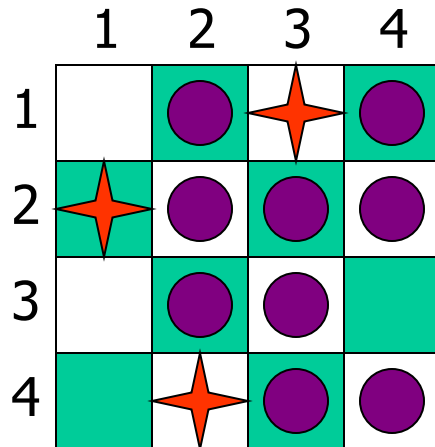
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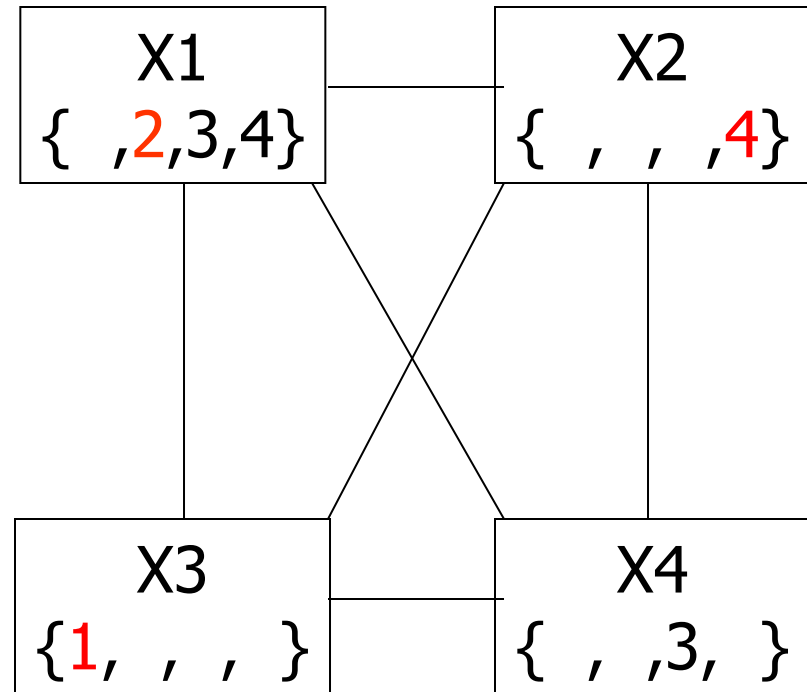
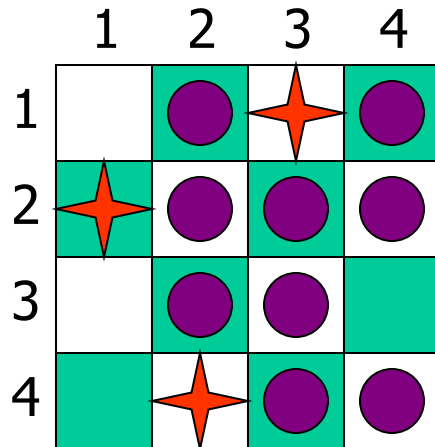
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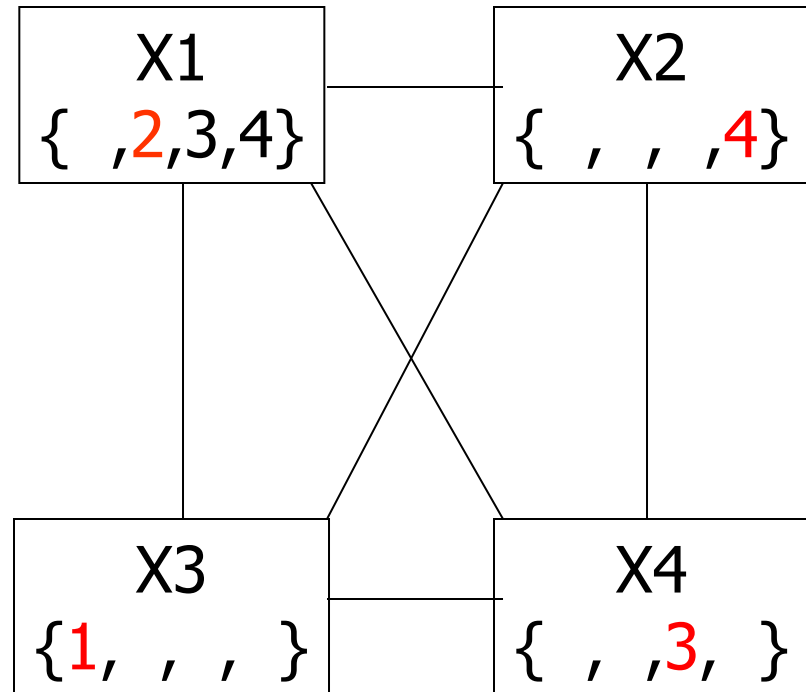
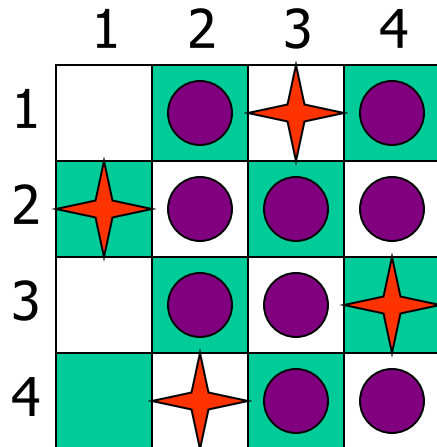
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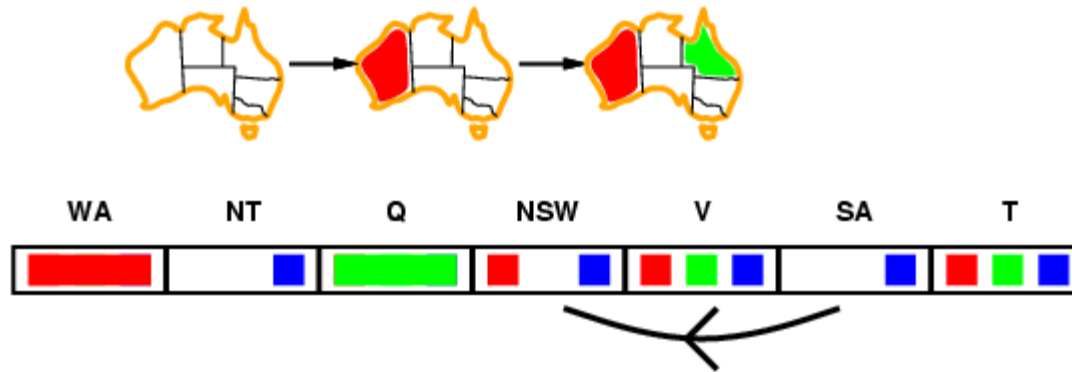


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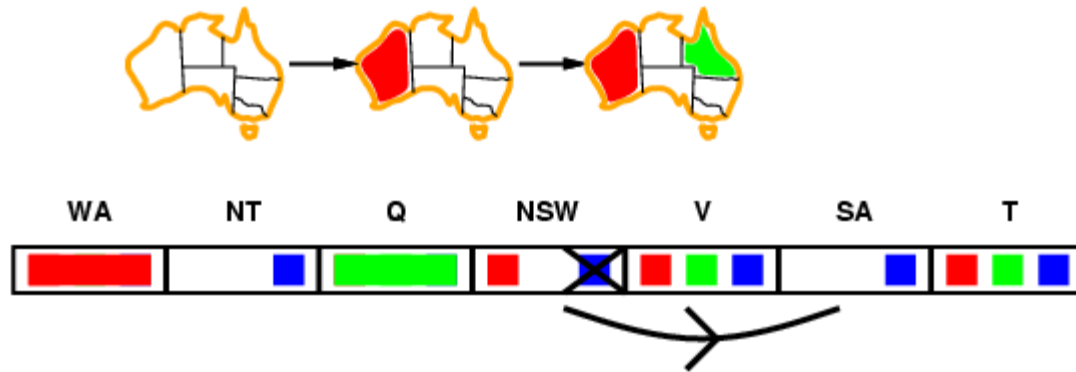
Constraint Propagation

- Simplest form of propagation makes each arc **consistent**
- Arc $X \rightarrow Y$ (link in constraint graph) is consistent iff
- for **every** value x of X there is **some** allowed y



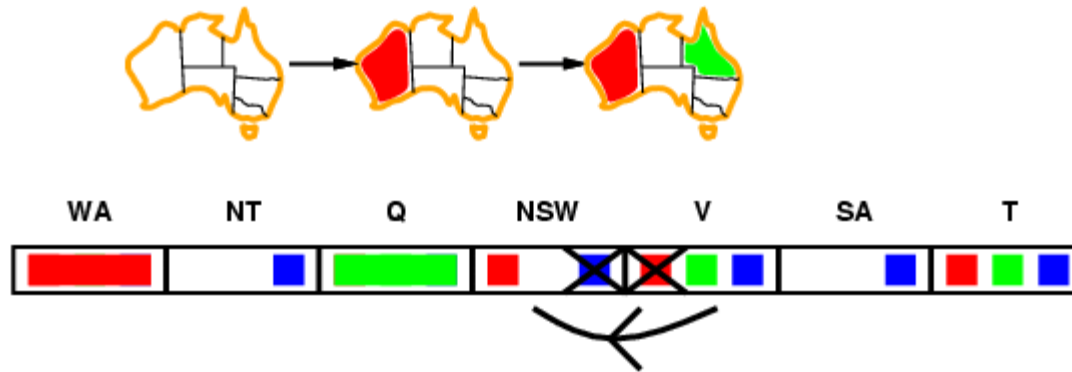
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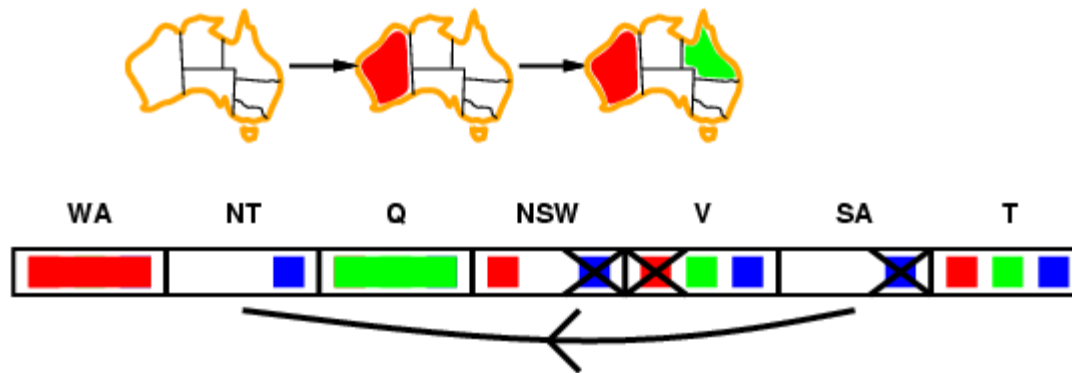
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- If X loses a value, neighbors of X need to be rechecked
- Arc consistency detects failure earlier than forward checking
- Can be run as a preprocessor or after each assignment

Arc Consistency Algorithm AC-3

function **AC-3**(*csp*) % returns the CSP, possibly with reduced domains

inputs: *csp*, a binary csp with variables $\{X_1, X_2, \dots, X_n\}$

local variables: *queue*, a queue of arcs initially the arcs in *csp*

while *queue* is not empty do

$(X_i, X_j) \leftarrow \text{REMOVE-FIRST}(\textit{queue})$

 if **REMOVE-INCONSISTENT-VALUES**(X_i, X_j) then

 for each X_k in **NEIGHBORS**[X_i] do

 add (X_k, X_i) to *queue*

function **REMOVE-INCONSISTENT-VALUES**(X_i, X_j) % returns *true* iff a value is removed

removed \leftarrow *false*

 for each x in **DOMAIN**[X_i] do

 if no value y in **DOMAIN**[X_j] allows (x,y) to satisfy the constraints between X_i and X_j

 then delete x from **DOMAIN**[X_i]; *removed* \leftarrow *true*

 return *removed*

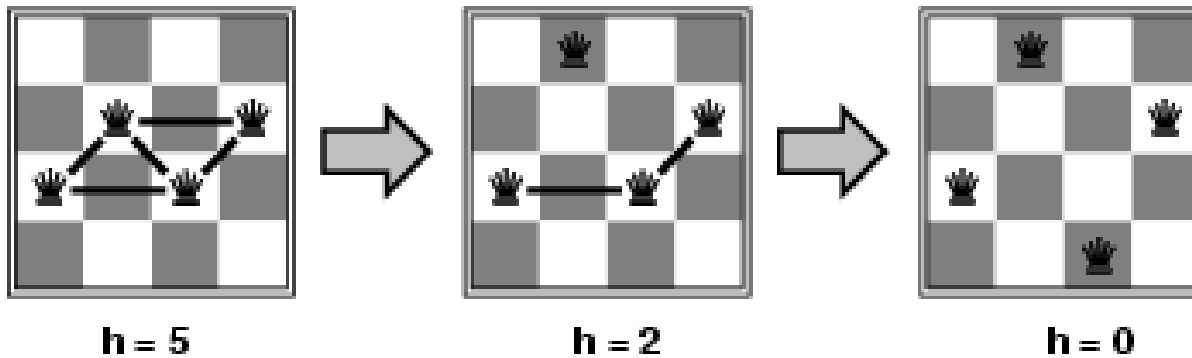
Time complexity: $O(n^2d^3)$

Local Search for CSPs

- **Hill-climbing methods typically work with "complete" states, i.e., all variables assigned**
- **To apply to CSPs:**
 - allow states with unsatisfied constraints
 - operators **reassign** variable values
- **Variable selection: randomly select any conflicted variable**
- **Value selection by **min-conflicts** heuristic:**
 - choose value that violates the fewest constraints
 - i.e., hill-climb with $h(n)$ = number of violated constraints

Example: n-queens

- **States:** 4 queens in 4 columns ($4^4 = 256$ states)
- **Actions:** move queen in column
- **Goal test:** no attacks
- **Evaluation:** $h(n) =$ number of attacks



- **Given random initial state, we can solve n -queens for large n with high probability**

Real-world CSPs

- **Assignment problems**
 - e.g., who teaches what class
- **Timetabling problems**
 - e.g., which class is offered when and where?
- **Transportation scheduling**
- **Factory scheduling**
- **Notice that many real-world problems involve real-valued variables**

Summary

- **CSPs are a special kind of problem:**
 - states defined by values of a fixed set of variables
 - goal test defined by constraints on variable values
- **Backtracking = depth-first search with one variable assigned per node**
- **Variable ordering and value selection heuristics help significantly**
- **Forward checking prevents assignments that guarantee later failure**
- **Constraint propagation (e.g., arc consistency) additionally constrains values and detects inconsistencies**
- **Iterative min-conflicts is usually effective in practice**