# Compositions of Tree-to-Tree Statistical Machine Translation Models

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Montréal — July 28, 2016

#### Required resource

• Parallel corpus

#### (containing sentences in both languages)

Additional resources	
<ul> <li>Language model data</li> </ul>	(text in the target language)
<ul> <li>Word alignments</li> </ul>	
• Parse trees	
• Various	(morphological analyzers, parsers)

# Statistical machine translation

Assume translation direction (English to Catalan)



but no parallel corpus for this pair



(similar motivation for non-deterministic pre- or post-processing steps)

#### Consequences

• 2 translation models

- (operating in sequence)
- Inefficiencies in sequential operation (due to sequential pruning)
- Theoretical guarantees missing for many operations

(e.g., tuning)

# Remedies [MayKniVog10] • Partial Evaluation [MayKniVog10] • Composition [MayKniVog10]

# Statistical machine translation

#### Vauquois triangle:



#### Translation model:

# Statistical machine translation

#### Vauquois triangle:



#### Translation model: tree-to-tree



- 2 Theoretical Model
- 3 Unweighted Compositions
- 4 Weighted Compositions

# Extended tree transducers

#### Definition

Transducer  $M = (Q, \Sigma, (q_1, q_2), R, wt)$  with

- finite set Q of states
- finite set  $\Sigma$  of symbols
- source initial state  $q_1$  and target initial state  $q_2$
- finite set R of rules
- weight assignment wt:  $R \rightarrow A$



(see below) (into a semiring)

# Extended tree transducers

#### Restrictions

Transducer  $M = (Q, \Sigma, (q_1, q_2), R, wt)$  is

- STSG if  $Q = \Sigma$  and state = root label
- SCFG if STSG and all rules are shallow
- $\varepsilon$ -free if no left-hand side is in Q
- strict if no right-hand side is in Q
- simple if  $\varepsilon$ -free and strict



 $\left(S, \underbrace{VBP}_{I}, NP\right) - \left(SENT, \underbrace{V}_{I}, NP\right)$ 

# Extended tree transducers

#### **Rules**:





#### Use in a derivation step:



#### Initial sentential form:



Final sentential form: No more linked states

#### Weight

- of a derivation: product of the weights of the used rules
- of a tree pair: sum of all derivations for the pair

#### Composition

Given two weighted relations  $\tau: T_{\Sigma} \times T_{\Delta} \to A$  and  $\tau': T_{\Delta} \times T_{\Gamma} \to A$ 

$$( au; au')(t,s) = \sum_{\upsilon\in \mathcal{T}_{\Delta}} au(t,\upsilon)\cdot au'(\upsilon,s)$$

# Background

#### 2 Theoretical Model

### Onweighted Compositions

#### Weighted Compositions

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#### $\mathsf{SCFG} \subsetneq \mathsf{STSG} \subsetneq \mathsf{TT}$

#### [Eis03] and folklore

Model	Composition closure	Reference
top-down transducer	1	[Eng75]
simple transducer	2	[ArnDau82]
other transducer	$\infty$	[EngFulMal16]

(top-down tree transducer = all left-hand sides shallow)

#### SCFG are closed under composition



# Second result

#### Difficulty for STSGs:



#### Extending the rules does not help:



. . .

Compositions of  $n \ge 2$  simple STSGs are as expressive as compositions of n simple transducers

Proof idea: Encode finite states in the intermediate tree(s)



#### Corollary

The composition closure for simple STSGs is obtained at level 2

#### Proof idea:

## $\mathsf{s}\text{-}\mathsf{STSG}\subsetneq\mathsf{s}\text{-}\mathsf{TT}\subsetneq\mathsf{s}\text{-}\mathsf{TT}^3=\mathsf{s}\text{-}\mathsf{STSG}^3=\mathsf{s}\text{-}\mathsf{STSG}^2$

Compositions of Tree-to-Tree Statistical Machine Translation Models

The composition hierarchy for the remaining STSGs is infinite

Proof idea: Inspect the corresponding proof for transducers and realize that the counterexamples can be generated by STSGs

# Background

2 Theoretical Model

3 Unweighted Compositions



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#### s-wTT; $wREL \subseteq s-wTT$

# $\mathsf{wREL}$ ; s- $\mathsf{wTT} \subseteq \mathsf{s-wTT}$

follows from [FulMalVog11] and [Kui99]

#### Lemma

#### $\mathsf{s}\text{-}\mathsf{w}\mathsf{T}\mathsf{T}\subseteq\mathsf{s}\mathsf{u}\text{-}\mathsf{T}\mathsf{T}_{\mathsf{inj}}\mathsf{ ; \mathsf{w}}\mathsf{R}\mathsf{E}\mathsf{L}$

Proof idea: Simply annotate rule identifier of applied rule to output and apply weights using the relabeling

 $(su-TT_{inj} = injective relations computed by simple unamb. transducers)$ 

The composition closure of weighted simple STSGs is obtained at level 2

#### Proof idea:

$$\begin{array}{ll} s\text{-}wTT\ ;\ s\text{-}wTT^2 &\subseteq su\text{-}TT_{inj}\ ;\ wREL\ ;\ s\text{-}wTT^2 \\ \subseteq\ su\text{-}TT_{inj}\ ;\ s\text{-}wTT^2 \subseteq\ su\text{-}TT_{inj}^2\ ;\ wREL\ ;\ s\text{-}wTT \\ \subseteq\ su\text{-}TT_{inj}^2\ ;\ s\text{-}wTT &\subseteq\ su\text{-}TT_{inj}^3\ ;\ wREL \\ \subseteq\ \underbrace{s\text{-}TT^2}_{injective}\ ;\ wREL &\subseteq\ su\text{-}TT_{inj}^2\ ;\ wREL \subseteq\ s\text{-}wTT^2 \end{array}$$

The composition hierarchy for the remaining weighted STSGs is infinite

# Proof idea: Utilize linking technique of [Mal15] to lift the corresponding unweighted result

Model	Composition closure
(weighted) SCFGs	1
(weighted) simple STSGs	2
(weighted) other STSGs	$\infty$

# Literature

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